

M58LR128KT M58LR128KB M58LR256KT M58LR256KB

128 or 256 Mbit (x16, multiple bank, multilevel interface, burst) 1.8 V supply flash memories

Features

- Supply voltage
 - V_{DD} = 1.7 V to 2.0 V for program, erase and read
 - V_{DDQ} = 1.7 V to 2.0 V for I/O buffers
 - V_{PP} = 9 V for fast program
- Synchronous/asynchronous read
 - Synchronous burst read mode: 54 MHz, 66 MHz
 - Asynchronous page read mode
 - Random access: 70 ns, 85 ns
- Synchronous burst read suspend
- Programming time
 - 2.5 µs typical word program time using Buffer Enhanced Factory Program command
- Memory organization
 - Multiple bank memory array: 8 Mbit banks for the M58LR128KT/B 16 Mbit banks for the M58LR256KT/B
 - Parameter blocks (top or bottom location)
- Dual operations
 - Program/erase in one bank while read in others
 - No delay between read and write operations
- Block locking
 - All blocks locked at power-up
 - Any combination of blocks can be locked with zero latency
 - WP for block lock-down
 - Absolute write protection with $V_{PP} = V_{SS}$



- 64 bit unique device number
- 2112 bit user programmable OTP cells
- Common flash interface (CFI)
- 100,000 program/erase cycles per block
- Electronic signature
 - Manufacturer code: 20h
 - Top device codes: M58LR128KT: 88C4h M58LR256KT: 880Dh
 - Bottom device codes M58LR128KB: 88C5h M58LR256KB: 880Eh
- The M58LR128KT/B is available in the ECOPACK-compliant VFBGA56 package.
- The M58LR256KT/B is available in the ECOPACK-compliant VFBGA79 and TFBGA88 packages.

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1 Description

The M58LR128KT/B and M58LR256KT/B are 128 Mbit (8 Mbit x16) and 256 Mbit (16 Mbit x 16) non-volatile Flash memories, respectively. They can be erased electrically at block level and programmed in-system on a word-by-word basis using a 1.7 V to 2.0 V V_{DD} supply for the circuitry and a 1.7 V to 2.0 V V_{DDQ} supply for the input/output pins. An optional 9 V V_{PP} power supply is provided to accelerate factory programming.

The devices feature an asymmetrical block architecture:

- The M58LR128KT/B have an array of 131 blocks, and are divided into 8 Mbit banks. There are 15 banks each containing 8 main blocks of 64 Kwords, and one parameter bank containing 4 parameter blocks of 16 Kwords and 7 main blocks of 64 Kwords.
- The M58LR256KT/B have an array of 259 blocks, and are divided into 16 Mbit banks. There are 15 banks each containing 16 main blocks of 64 Kwords, and one parameter bank containing 4 parameter blocks of 16 Kwords and 15 main blocks of 64 Kwords.

The multiple bank architecture allows dual operations. While programming or erasing in one bank, read operations are possible in other banks. Only one bank at a time is allowed to be in program or erase mode. It is possible to perform burst reads that cross bank boundaries. The bank architecture is summarized in *Table 2*, and the memory map is shown in *Figure 4*. The parameter blocks are located at the top of the memory address space for the M58LR128KT and M58LR256KT, and at the bottom for the M58LR128KB and M58LR256KB.

Each block can be erased separately. Erase can be suspended to perform a program or read operation in any other block, and then resumed. Program can be suspended to read data at any memory location except for the one being programmed, and then resumed. Each block can be programmed and erased over 100 000 cycles using the supply voltage V_{DD} . There is a buffer enhanced factory programming command available to speed up programming.

Program and erase commands are written to the command interface of the memory. An internal Program/Erase Controller manages the timings necessary for program and erase operations. The end of a program or erase operation can be detected and any error conditions identified in the Status Register. The command set required to control the memory is consistent with JEDEC standards.

The device supports synchronous burst read and asynchronous read from all blocks of the memory array; at power-up the device is configured for asynchronous read. In synchronous burst read mode, data is output on each clock cycle at frequencies of up to 66 MHz. The synchronous burst read operation can be suspended and resumed.

The device features an automatic standby mode. When the bus is inactive during asynchronous read operations, the device automatically switches to automatic standby mode. In this condition the power consumption is reduced to the standby value and the outputs are still driven.

The M58LRxxxKT/B features an instant, individual block locking scheme that allows any block to be locked or unlocked with no latency, enabling instant code and data protection. All blocks have three levels of protection. They can be locked and locked-down individually preventing any accidental programming or erasure. There is an additional hardware protection against program and erase. When $V_{PP} \leq V_{PPLK}$ all blocks are protected against program or erase. All blocks are locked at power-up.



The device includes 17 Protection Registers and 2 Protection Register locks, one for the first Protection Register and the other for the 16 one-time-programmable (OTP) Protection Registers of 128 bits each. The first Protection Register is divided into two segments: a 64 bit segment containing a unique device number written by Numonyx, and a 64 bit segment OTP by the user. The user programmable segment can be permanently protected. *Figure 6*, shows the Protection Register memory map.

The M58LR128KT/B is offered in a VFBGA56 7.7 x 9 mm, 0.75 mm package, and the M58LR256KT/B is offered in a VFBGA79 9 x 11mm, 0.75 mm package and TFGBA88 8 x 10mm, 0.8 mm package.

All devices are supplied with all the bits erased (set to '1').



Figure 1. Logic diagram

1. Amax is equal to A22 in the M58LR128KT/B and, to A23 in the M58LR256KT/B.



Signal name	Function	Direction	
A0-Amax ⁽¹⁾	Address inputs	Inputs	
DQ0-DQ15	Data input/outputs, command inputs	I/O	
Ē	Chip Enable	Input	
G	Output Enable	Input	
W	Write Enable	Input	
RP	Reset	Input	
WP	Write Protect	Input	
К	Clock	Input	
Ē	Latch Enable	Input	
WAIT	Wait	Output	
V _{DD}	Supply voltage		
V _{DDQ}	Supply voltage for input/output buffers		
V _{PP}	Optional supply voltage for fast program & erase		
V _{SS}	Ground		
V _{SSQ}	Ground input/output supply		

Table 1.Signal names

1. Amax is equal to A22 in the M58LR128KT/B and, to A23 in the M58LR256KT/B.

1 2 3 7 8 4 5 6 A8 А A11 A18 A6 A4 V_{SS} V_{DD} $\overline{\mathsf{RP}}$ κ в A9 A20 A5 A12 A17 AЗ L $\overline{\mathsf{W}}$ С A10 A21 A19 Α7 A2 A13 D WP A14 WAIT A16 DQ12 A22 A15 A1 DQ6 DQ1 DQ15 Ē Е DQ4 DQ2 A0 V_{DDQ} G F DQ10 DQ0 DQ14 DQ13 DQ11 DQ9 Vss G DQ7 DQ5 DQ3 VDDQ DQ8 V_{DD} VSSQ V_{SSQ}



1. This package is available only for the M58LR128KT/B devices.

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	1	2	3	4	5	6	7	8
A		DU					DU	DU
в	A4	A18	A19	V _{SS}	V _{DD}	NC	A21	(A11
с	A5	NC	A23	V _{ss}	NC	(к)	A22	A12
D	A3	A17	NC	V _{PP}	NC	NC	A9	A13
E	A2	A7	NC	WP	([])	A20	A10	A15
F	(A1)	A6	NC	RP	Ŵ	A8	A14	A16
G	AO	DQ8	DQ2	DQ10	DQ5	(DQ13)	WAIT	NC
н	NC	DQ0	DQ1	DQ3	DQ12	DQ14	DQ7	NC
J	NC	Ģ	DQ9	(DQ11)	DQ4	DQ6	(DQ15)	V _{DDQ}
к	Ē	NC	NC	NC	NC	NC	V _{DDQ}	NC
L	V _{SS}	V _{SS}	V _{DDQ}	V _{DD}	V _{SS}	V _{SS}	V _{SS}	V _{SS}
М	(DU)	DU					DU	(DU)

Figure 3. TFBGA88 connections (top view through package)

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Figure 1: 7x9 Active-Ball Matrix for 256-Mbit Density in VF BGA Package



N numonyx

Number	Bank size	Parameter blocks	Main blocks
Parameter Bank	8 Mbits	4 blocks of 16 Kwords	7 blocks of 64 Kwords
Bank 1	8 Mbits	-	8 blocks of 64 Kwords
Bank 2	8 Mbits	-	8 blocks of 64 Kwords
Bank 3	8 Mbits	-	8 blocks of 64 Kwords
Bank 14	8 Mbits	-	8 blocks of 64 Kwords
Bank 15	8 Mbits	-	8 blocks of 64 Kwords

Table 2. M58LR128KT/B bank architecture

Figure 4. M58LR128KT/B memory map



Number	Bank size	Parameter blocks	Main blocks
Parameter bank	16 Mbits	4 blocks of 16 Kwords	15 blocks of 64 Kwords
Bank 1	16 Mbits	-	16 blocks of 64 Kwords
Bank 2	16 Mbits	-	16 blocks of 64 Kwords
Bank 3	16 Mbits	-	16 blocks of 64 Kwords
Bank 14	16 Mbits	-	16 blocks of 64 Kwords
Bank 15	16 Mbits	-	16 blocks of 64 Kwords

Table 3. M58LR256KT/B bank architecture





2 Signal descriptions

See *Figure 1: Logic diagram* and *Table 1: Signal names* for a brief overview of the signals connected to this device.

2.1 Address inputs (A0-Amax)

Amax is the highest order address input. It is equal to A22 in the M58LR128KT/B and, to A23 in the M58LR256KT/B. The address inputs select the cells in the memory array to access during bus read operations. During bus write operations they control the commands sent to the command interface of the Program/Erase Controller.

2.2 Data inputs/outputs (DQ0-DQ15)

The data I/O output the data stored at the selected address during a bus read operation or input a command or the data to be programmed during a bus write operation.

2.3 Chip Enable (\overline{E})

The Chip Enable input activates the memory control logic, input buffers, decoders and sense amplifiers. When Chip Enable is at V_{IL} and Reset is at V_{IH} the device is in active mode. When Chip Enable is at V_{IH} the memory is deselected, the outputs are high impedance and the power consumption is reduced to the standby level.

2.4 Output Enable (G)

The Output Enable input controls data outputs during the bus read operation of the memory.

2.5 Write Enable (\overline{W})

The Write Enable input controls the bus write operation of the memory's command interface. The data and address inputs are latched on the rising edge of Chip Enable or Write Enable, whichever occurs first.

2.6 Write Protect (WP)

Write Protect is an input that gives an additional hardware protection for each block. When Write Protect is at V_{IL} , the lock-down is enabled and the protection status of the locked-down blocks cannot be changed. When Write Protect is at V_{IH} , the lock-down is disabled and the locked-down blocks can be locked or unlocked. (refer to *Table 17: Lock status*).

2.7 Reset (\overline{RP})

The Reset input provides a hardware reset of the memory. When Reset is at V_{IL} , the memory is in reset mode: the outputs are high impedance and the current consumption is reduced to the Reset supply current I_{DD2} . Refer to *Table 22: DC characteristics - currents* for the value of I_{DD2} . After Reset all blocks are in the locked state and the Configuration Register is reset. When Reset is at V_{IH} , the device is in normal operation. Exiting reset mode the device enters asynchronous read mode, and a negative transition of Chip Enable or Latch Enable is required to ensure valid data outputs.

2.8 Latch Enable (\overline{L})

Latch Enable latches the address bits on its rising edge. The address latch is transparent when Latch Enable is at V_{IL} and it is inhibited when Latch Enable is at V_{IH} .

2.9 Clock (K)

The Clock input synchronizes the memory to the microcontroller during synchronous read operations; the address is latched on a Clock edge (rising or falling, according to the configuration settings) when Latch Enable is at V_{IL} . Clock is ignored during asynchronous read and in write operations.

2.10 Wait (WAIT)

Wait is an output signal used during synchronous read to indicate whether the data on the output bus are valid. This output is high impedance when Chip Enable is at V_{IH} , Output Enable is at V_{IH} or Reset is at V_{IL} . It can be configured to be active during the wait cycle or one clock cycle in advance.

2.11 V_{DD} supply voltage

 V_{DD} provides the power supply to the internal core of the memory device. It is the main power supply for all operations (read, program and erase).

2.12 V_{DDQ} supply voltage

 V_{DDQ} provides the power supply to the I/O pins and enables all outputs to be powered independently from $V_{DD}.$ V_{DDQ} can be tied to V_{DD} or can use a separate supply.



2.13 V_{PP} program supply voltage

 V_{PP} is both a control input and a power supply pin. The two functions are selected by the voltage range applied to the pin.

If V_{PP} is kept in a low voltage range (0V to V_{DDQ}) V_{PP} is seen as a control input. In this case a voltage lower than V_{PPLK} provides absolute protection against program or erase, while V_{PP} in the V_{PP1} range enables these functions (see Tables 22 and 23, DC characteristics for the relevant values). V_{PP} is only sampled at the beginning of a program or erase. A change in its value after the operation has started does not have any effect and program or erase operations continue.

If V_{PP} is in the range of V_{PPH} it acts as a power supply pin. In this condition V_{PP} must be stable until the program/erase algorithm is completed.

2.14 V_{SS} ground

V_{SS} ground is the reference for the core supply. It must be connected to the system ground.

2.15 V_{SSQ} ground

 V_{SSQ} ground is the reference for the input/output circuitry driven by $V_{DDQ}.~V_{SSQ}$ must be connected to V_{SS}

Note: Each device in a system should have V_{DD} , V_{DDQ} and V_{PP} decoupled with a 0.1 μ F ceramic capacitor close to the pin (high frequency, inherently low inductance capacitors should be as close as possible to the package). See Figure 10: AC measurement load circuit. The PCB track widths should be sufficient to carry the required V_{PP} program and erase currents.

3 Bus operations

There are six standard bus operations that control the device. These are bus read, bus write, address latch, output disable, standby and reset. See *Table 4: Bus operations* for a summary.

Typically glitches of less than 5 ns on Chip Enable or Write Enable are ignored by the memory and do not affect bus write operations.

3.1 Bus read

Bus read operations are used to output the contents of the memory array, the electronic signature, the Status Register and the Common Flash interface. Both Chip Enable and Output Enable must be at V_{IL} to perform a read operation. The Chip Enable input should be used to enable the device, and Output Enable should be used to gate data onto the output. The data read depends on the previous command written to the memory (see command interface section). See Figures *11*, *12* and *13* read AC waveforms, and Tables *24* and *25* read AC characteristics for details of when the output becomes valid.

3.2 Bus write

Bus write operations write commands to the memory or latch input data to be programmed. A bus write operation is initiated when Chip Enable and Write Enable are at V_{IL} with Output Enable at V_{IH} . Commands, input data and addresses are latched on the rising edge of Write Enable or Chip Enable, whichever occurs first. The addresses must be latched prior to the write operation by toggling Latch Enable (when Chip Enable is at V_{IL}). The Latch Enable must be tied to V_{IH} during the bus write operation.

See Figures 17 and 18, write AC waveforms, and Tables 26 and 27, write AC characteristics for details of the timing requirements.

3.3 Address Latch

Address latch operations input valid addresses. Both Chip enable and Latch Enable must be at V_{IL} during address latch operations. The addresses are latched on the rising edge of Latch Enable.

3.4 Output disable

The outputs are high impedance when the Output Enable is at V_{IH}.



3.5 Standby

Standby disables most of the internal circuitry allowing a substantial reduction of the current consumption. The memory is in standby when Chip Enable and Reset are at V_{IH}. The power consumption is reduced to the standby level I_{DD3} and the outputs are set to high impedance, independently from the Output Enable or Write Enable inputs. If Chip Enable switches to V_{IH} during a program or erase operation, the device enters standby mode when finished.

3.6 Reset

During reset mode the memory is deselected and the outputs are high impedance. The memory is in reset mode when Reset is at V_{IL} . The power consumption is reduced to the reset level, independently from the Chip Enable, Output Enable or Write Enable inputs. If Reset is pulled to V_{SS} during a program or erase, this operation is aborted and the memory content is no longer valid.

Operation	Ē	G	W	Ē	RP	WAIT ⁽²⁾	DQ15-DQ0
Bus read	V_{IL}	V_{IL}	V _{IH}	V _{IL} ⁽³⁾	V _{IH}		Data output
Bus write	V_{IL}	V _{IH}	V _{IL}	V _{IL} ⁽³⁾	V _{IH}		Data input
Address latch	V_{IL}	Х	V _{IH}	V _{IL}	V _{IH}		Data output or Hi-Z ⁽⁴⁾
Output disable	V_{IL}	V_{IH}	V _{IH}	Х	V _{IH}	Hi-Z	Hi-Z
Standby	V _{IH}	Х	Х	Х	V _{IH}	Hi-Z	Hi-Z
Reset	Х	Х	Х	Х	V _{IL}	Hi-Z	Hi-Z

Table 4.Bus operations⁽¹⁾

1. X = 'don't care'

2. WAIT signal polarity is configured using the Set Configuration Register command.

3. \overline{L} can be tied to V_{IH} if the valid address has been previously latched.

4. Depends on \overline{G} .

4 Command interface

All bus write operations to the memory are interpreted by the command interface. Commands consist of one or more sequential bus write operations. An internal Program/Erase Controller manages all timings and verifies the correct execution of the program and erase commands. The Program/Erase Controller provides a Status Register whose output may be read at any time to monitor the progress or the result of the operation.

The command interface is reset to read mode when power is first applied, when exiting from Reset or whenever V_{DD} is lower than V_{LKO} . Command sequences must be followed exactly. Any invalid combination of commands are ignored.

Refer to *Table 5: Command codes*, *Table 6: Standard commands*, *Table 7: Factory commands* and *Appendix D: Command interface state tables* for a summary of the command interface.

Hex Code	Command
01h	Block Lock Confirm
03h	Set Configuration Register Confirm
10h	Alternative Program Setup
20h	Block Erase Setup
2Fh	Block Lock-Down Confirm
40h	Program Setup
50h	Clear Status Register
60h	Block Lock Setup, Block Unlock Setup, Block Lock Down Setup and Set Configuration Register Setup
70h	Read Status Register
80h	Buffer Enhanced Factory Program Setup
90h	Read Electronic Signature
98h	Read CFI Query
B0h	Program/Erase Suspend
BCh	Blank Check Setup
C0h	Protection Register Program
CBh	Blank Check Confirm
D0h	Program/Erase Resume, Block Erase Confirm, Block Unlock Confirm, Buffer Program or Buffer Enhanced Factory Program Confirm
E8h	Buffer Program
FFh	Read Array

Table 5. Command codes



4.1 Read Array command

The Read Array command returns the addressed bank to read array mode.

One bus write cycle is required to issue the Read Array command. Once a bank is in read array mode, subsequent read operations outputs the data from the memory array.

A Read Array command can be issued to any banks while programming or erasing in another bank.

If the Read Array command is issued to a bank currently executing a program or erase operation, the bank returns to read array mode but the program or erase operation continues, however the data output from the bank is not guaranteed until the program or erase operation has finished. The read modes of other banks are not affected.

4.2 Read Status Register command

The device contains a Status Register that monitors program or erase operations.

The Read Status Register command reads the contents of the Status Register for the addressed bank.

One bus write cycle is required to issue the Read Status Register command. Once a bank is in read Status Register mode, subsequent read operations output the contents of the Status Register.

The Status Register data is latched on the falling edge of the Chip Enable or Output Enable signals. Either Chip Enable or Output Enable must be toggled to update the Status Register data.

The Read Status Register command can be issued at any time, even during program or erase operations. The Read Status Register command only changes the read mode of the addressed bank. The read modes of other banks are not affected. only asynchronous read and single synchronous read operations should be used to read the Status Register. A Read Array command is required to return the bank to read array mode.

See *Table 10* for the description of the Status Register bits.

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4.3 Read Electronic Signature command

The Read Electronic Signature command reads the manufacturer and device codes, the lock status of the addressed bank, the Protection Register, and the Configuration Register.

One bus write cycle is required to issue the Read Electronic Signature command. Once a bank is in read electronic signature mode, subsequent read operations in the same bank output the manufacturer code, the device code, the lock status of the addressed bank, the Protection Register, or the Configuration Register (see *Table 9*).

The Read Electronic Signature command can be issued at any time, even during program or erase operations, except during Protection Register Program operations. Dual operations between the parameter bank and the electronic signature location are not allowed (see *Table 16: Dual operation limitations* for details).

If a Read Electronic Signature command is issued to a bank that is executing a program or erase operation, the bank goes into read electronic signature mode. Subsequent bus read cycles output the electronic signature data and the Program/Erase Controller continue to program or erase in the background.

The Read Electronic Signature command only changes the read mode of the addressed bank. The read modes of other banks are not affected. Only asynchronous read and single synchronous read operations should be used to read the electronic signature. A Read Array command is required to return the bank to read array mode.

4.4 Read CFI Query command

The Read CFI Query command reads data from the common Flash interface (CFI).

One bus write cycle is required to issue the Read CFI Query command. Once a bank is in read CFI query mode, subsequent bus read operations in the same bank read from the common Flash interface.

The Read CFI Query command can be issued at any time, even during program or erase operations.

If a Read CFI Query command is issued to a bank that is executing a program or erase operation the bank gos into read CFI query mode. Subsequent bus read cycles output the CFI data and the Program/Erase Controller continues to program or erase in the background.

The Read CFI Query command only changes the read mode of the addressed bank. The read modes of other banks are not affected. Only asynchronous read and single synchronous read operations should be used to read from the CFI. A Read Array command is required to return the bank to read array mode. Dual operations between the parameter bank and the CFI memory space are not allowed (see *Table 16: Dual operation limitations* for details).

See *Appendix B: Common Flash interface* and Tables *45*, *46*, *47*, *48*, *49*, *50*, *51*, *52*, *53* and *54* for details on the information contained in the common Flash interface memory area.



4.5 Clear Status Register command

The Clear Status Register command resets (set to '0') all error bits (SR1, 3, 4 and 5) in the Status Register.

One bus write cycle is required to issue the Clear Status Register command. The Clear Status Register command does not affect the read mode of the bank.

The error bits in the Status Register do not automatically return to '0' when a new command is issued. The error bits in the Status Register should be cleared before attempting a new program or erase command.

4.6 Block Erase command

The Block Erase command erases a block. It sets all the bits within the selected block to '1'. All previous data in the block is lost.

If the block is protected then the erase operation aborts, the data in the block does not change and the Status Register outputs the error.

Two bus write cycles are required to issue the command.

- The first bus cycle sets up the Block Erase command.
- The second latches the block address and starts the Program/Erase Controller.

If the second bus cycle is not the Block Erase Confirm code, Status Register bits SR4 and SR5 are set and the command is aborted.

Once the command is issued the bank enters Read Status Register mode and any read operation within the addressed bank outputs the contents of the Status Register. A Read Array command is required to return the bank to read array mode.

During block erase operations the bank containing the block being erased only accepts the Read Array, Read Status Register, Read Electronic Signature, Read CFI Query and the Program/Erase Suspend commands; all other commands are ignored.

The block erase operation aborts if Reset, \overline{RP} , goes to V_{IL}. As data integrity cannot be guaranteed when the block erase operation is aborted, the block must be erased again.

Refer to *Section 8: Dual operations and multiple bank architecture* for detailed information about simultaneous operations allowed in banks not being erased.

Typical erase times are given in Table 18: Program/erase times and endurance cycles.

See *Appendix C*, *Figure 28: Block erase flowchart and pseudocode* for a suggested flowchart for using the Block Erase command.



4.7 Blank Check command

The Blank Check command checks whether a main array block has been completely erased. Only one block at a time can be checked. To use the Blank Check command V_{PP} must be equal to V_{PPH} . If V_{PP} is not equal to V_{PPH} , the device ignores the command and no error is shown in the Status Register.

Two bus cycles are required to issue the Blank Check command:

- The first bus cycle writes the Blank Check command (BCh) to any address in the block to be checked.
- The second bus cycle writes the Blank Check Confirm command (CBh) to any address in the block to be checked and starts the blank check operation.

If the second bus cycle is not Blank Check Confirm, Status Register bits SR4 and SR5 are set to '1' and the command aborts.

Once the command is issued, the addressed bank automatically enters the Status Register mode and further reads within the bank output the Status Register contents.

The only operation permitted during blank check is Read Status Register. Dual operations are not supported while a blank check operation is in progress. Blank check operations cannot be suspended and are not allowed while the device is in program/erase suspend.

The SR7 Status Register bit indicates the status of the blank check operation in progress. SR7 = '0' means that the Blank Check operation is still ongoing, and SR7 = '1' means that the operation is complete.

The SR5 Status Register bit goes High (SR5 = '1') to indicate that the blank check operation has failed.

At the end of the operation the bank remains in the read Status Register mode until another command is written to the command interface.

See *Appendix C*, *Figure 25: Blank check flowchart and pseudocode* for a suggested flowchart for using the Blank Check command.

Typical blank check times are given in *Table 18: Program/erase times and endurance cycles*.

4.8 **Program command**

The Program command programs a single word to the memory array.

If the block being programmed is protected, then the program operation aborts, the data in the block does not change and the Status Register outputs the error.

Two bus write cycles are required to issue the Program command:

- The first bus cycle sets up the Program command.
- The second latches the address and data to be programmed and starts the Program/Erase Controller.

Once the programming has started, read operations in the bank being programmed output the Status Register content.

During a program operation, the bank containing the word being programmed only accepts the Read Array, Read Status Register, Read Electronic Signature, Read CFI Query and the Program/Erase Suspend commands; all other commands are ignored. A Read Array command is required to return the bank to read array mode.

Refer to *Section 8: Dual operations and multiple bank architecture* for detailed information about simultaneous operations allowed in banks not being programmed.

Typical program times are given in Table 18: Program/erase times and endurance cycles.

The program operation aborts if Reset, \overline{RP} , goes to V_{IL}. As data integrity cannot be guaranteed when the program operation is aborted, the word must be reprogrammed.

See *Appendix C*, *Figure 24: Program flowchart and pseudocode* for the flowchart for using the Program command.

4.9 Buffer Program command

The Buffer Program Command uses the device's 32-word write buffer to speed up programming. Up to 32 words can be loaded into the write buffer. The Buffer Program command dramatically reduces in-system programming time compared to the standard non-buffered Program command.

Four successive steps are required to issue the Buffer Program command:

1. The first bus write cycle sets up the Buffer Program command. The setup code can be addressed to any location within the targeted block.

After the first bus write cycle, read operations in the bank output the contents of the Status Register. Status Register bit SR7 should be read to check that the buffer is available (SR7 = 1). If the buffer is not available (SR7 = 0), re-issue the Buffer Program command to update the Status Register contents.

 The second bus write cycle sets up the number of words to be programmed. Value n is written to the same block address, where n+1 is the number of words to be programmed.

- 3. Use n+1 bus write cycles to load the address and data for each word into the write buffer. Addresses must lie within the range from the start address to the start address + n, where the start address is the location of the first data to be programmed. Optimum performance is obtained when the start address corresponds to a 32 word boundary.
- 4. The final bus write cycle confirms the Buffer Program command and starts the program operation.

All the addresses used in the Buffer Program operation must lie within the same block.

Invalid address combinations or failing to follow the correct sequence of bus write cycles sets an error in the Status Register and aborts the operation without affecting the data in the memory array.

If the Status Register bits SR4 and SR5 are set to '1', the Buffer Program command is not accepted. Clear the Status Register before re-issuing the command.

If the block being programmed is protected an error sets in the Status Register and the operation aborts without affecting the data in the memory array.

During Buffer Program operations the bank being programmed only accepts the Read Array, Read Status Register, Read Electronic Signature, Read CFI Query and the Program/Erase Suspend commands; all other commands are ignored.

Refer to *Section 8: Dual operations and multiple bank architecture* for detailed information about simultaneous operations allowed in banks not being programmed.

See *Appendix C*, *Figure 26: Buffer program flowchart and pseudocode* for a suggested flowchart on using the Buffer Program command.

4.10 Buffer Enhanced Factory Program command

The Buffer Enhanced Factory Program command has been specially developed to speed up programming in manufacturing environments where the programming time is critical.

It is used to program one or more write buffer(s) of 32 words to a block. Once the device enters Buffer Enhanced Factory Program mode, the write buffer can be reloaded any number of times as long as the address remains within the same block. Only one block can be programmed at a time.

If the block being programmed is protected, then the Program operation aborts the data in the block does not change, and the Status Register outputs the error.

The use of the Buffer Enhanced Factory Program command requires the following operating conditions:

- V_{PP} must be set to V_{PPH}
- V_{DD} must be within operating range
- Ambient temperature T_A must be 30°C ± 10°C
- The targeted block must be unlocked
- The start address must be aligned with the start of a 32 word buffer boundary
- The address must remain the start address throughout programming.

Dual operations are not supported during the Buffer Enhanced Factory Program operation and the command cannot be suspended.



The Buffer Enhanced Factory Program Command consists of three phases: the setup phase, the program and verify phase, and the exit phase, Please refer to *Table 7: Factory commands* for detailed information.

4.10.1 Setup phase

The Buffer Enhanced Factory Program command requires two bus write cycles to initiate the command:

- The first bus write cycle sets up the Buffer Enhanced Factory Program command.
- The second bus write cycle confirms the command.

After the confirm command is issued, read operations output the contents of the Status Register. The Read Status Register command must not be issued or it is interpreted as data to program.

The Status Register P/EC bit SR7 should be read to check that the P/EC is ready to proceed to the next phase.

If an error is detected, SR4 goes high (set to '1') and the Buffer Enhanced Factory Program operation is terminated. See *Section 5: Status Register* for details on the error.

4.10.2 Program and verify phase

The program and verify phase requires 32 cycles to program the 32 words to the write buffer. The data is stored sequentially, starting at the first address of the write buffer, until the write buffer is full (32 words). To program less than 32 words, the remaining words should be programmed with FFFFh.

Three successive steps are required to issue and execute the program and verify phase of the command:

- 1. Use one bus write operation to latch the start address and the first word to be programmed. The Status Register bank write status bit SR0 should be read to check that the P/EC is ready for the next word.
- 2. Each subsequent word to be programmed is latched with a new bus write operation. The address must remain the start address as the P/EC increments the address location. If any address is given that is not in the same block as the start address, the program and verify phase terminates. Status Register bit SR0 should be read between each bus write cycle to check that the P/EC is ready for the next word.
- 3. Once the write buffer is full, the data is programmed sequentially to the memory array. After the program operation the device automatically verifies the data and reprograms if necessary.

The program and verify phase can be repeated, without re-issuing the command, to program additional 32 word locations as long as the address remains in the same block.

4. Finally, after all words, or the entire block have been programmed, write one bus write operation to any address outside the block containing the start address, to terminate program and verify phase.

Status Register bit SR0 must be checked to determine whether the program operation is finished. The Status Register may be checked for errors at any time but it must be checked after the entire block has been programmed.

4.10.3 Exit phase

Status Register P/EC bit SR7 set to '1' indicates that the device has exited the buffer enhanced factory program operation and returned to read Status Register mode. A full Status Register check should be done to ensure that the block has been successfully programmed. See *Section 5: Status Register* for more details.

For optimum performance the Buffer Enhanced Factory Program command should be limited to a maximum of 100 program/erase cycles per block. If this limit is exceeded the internal algorithm continues to work properly but some degradation in performance is possible. Typical program times are given in *Table 18*.

See Appendix C, Figure 32: Buffer enhanced factory program flowchart and pseudocode for a suggested flowchart on using the Buffer Enhanced Factory Program command.

4.11 Program/Erase Suspend command

The Program/Erase Suspend command pauses a program or block erase operation. The command can be addressed to any bank.

The Program/Erase Resume command is required to restart the suspended operation.

One bus write cycle is required to issue the Program/Erase Suspend command. Once the Program/Erase Controller has paused bits SR7, SR6 and/ or SR2 of the Status Register are set to '1'.

The following commands are accepted during Program/Erase Suspend:

- Program/Erase Resume
- Read Array (data from erase-suspended block or program-suspended word is not valid)
- Read Status Register
- Read Electronic Signature
- Read CFI Query
- Clear Status Register

Additionally, if the suspended operation was a block erase then the following commands are also accepted:

- Set Configuration Register
- Program (except in erase-suspended block)
- Buffer Program (except in erase suspended blocks)
- Block Lock
- Block Lock-Down
- Block Unlock.

During an erase suspend the block being erased can be protected by issuing the Block Lock or Block Lock-Down commands. When the Program/Erase Resume command is issued the operation completes.

It is possible to accumulate multiple suspend operations. For example, it is possible to suspend an erase operation, start a program operation, suspend the program operation, and then read the array.



If a Program command is issued during a block erase suspend, the erase operation cannot be resumed until the program operation has completed.

The Program/Erase Suspend command does not change the read mode of the banks. If the suspended bank was in read Status Register, read electronic signature or read CFI query mode, the bank remains in that mode and outputs the corresponding data.

Refer to *Section 8: Dual operations and multiple bank architecture* for detailed information about simultaneous operations allowed during Program/Erase Suspend.

During a Program/Erase Suspend, the device can be placed in standby mode by taking Chip Enable to V_{IH} . Program/erase is aborted if Reset, \overline{RP} , goes to V_{II} .

See Appendix C, Figure 27: Program suspend and resume flowchart and pseudocode, and Figure 29: Erase suspend and resume flowchart and pseudocode for flowcharts for using the Program/Erase Suspend command.

4.12 Program/Erase Resume command

The Program/Erase Resume command restarts the program or erase operation suspended by the Program/Erase Suspend command. One bus write cycle is required to issue the command. The command can be issued to any address.

The Program/Erase Resume command does not change the read mode of the banks. If the suspended bank was in read Status Register, read electronic signature or read CFI query mode the bank remains in that mode and outputs the corresponding data.

If a program command is issued during a block erase suspend, then the erase cannot be resumed until the program operation has completed.

See Appendix C, Figure 27: Program suspend and resume flowchart and pseudocode, and Figure 29: Erase suspend and resume flowchart and pseudocode for flowcharts for using the Program/Erase Resume command.

4.13 Protection Register Program command

The Protection Register Program command programs the user OTP segments of the Protection Register and the two Protection Register locks.

The device features 16 OTP segments of 128 bits and one OTP segment of 64 bits, as shown in *Figure 6: Protection Register memory map*.

The segments are programmed one word at a time. When shipped all bits in the segment are set to '1'. The user can only program the bits to '0'.

Two bus write cycles are required to issue the Protection Register Program command.

- The first bus cycle sets up the Protection Register Program command.
- The second latches the address and data to be programmed to the Protection Register and starts the Program/Erase Controller.

Read operations to the bank being programmed output the Status Register content after the program operation has started.

Attempting to program a previously protected Protection Register results in a Status Register error.

The Protection Register Program cannot be suspended. Dual operations between the parameter bank and the Protection Register memory space are not allowed (see *Table 16: Dual operation limitations* for details)

The two Protection Register locks protect the OTP segments from further modification. The protection of the OTP segments is not reversible. Refer to *Figure 6: Protection Register memory map* and *Table 9: Protection Register locks* for details on the lock bits.

See *Appendix C*, *Figure 31: Protection Register program flowchart and pseudocode* for a flowchart for using the Protection Register Program command.

4.14 Set Configuration Register command

The Set Configuration Register command writes a new value to the Configuration Register.

Two bus write cycles are required to issue the Set Configuration Register comman:

- The first cycle sets up the Set Configuration Register command and the address corresponding to the Configuration Register content.
- The second cycle writes the Configuration Register data and the confirm command.

The Configuration Register data must be written as an address during the bus write cycles, that is A0 = CR0, A1 = CR1, ..., A15 = CR15. Addresses A16-Amax are ignored.

Read operations output the array content after the Set Configuration Register command is issued.

The Read Electronic Signature command is required to read the updated contents of the Configuration Register.

4.15 Block Lock command

The Block Lock command locks a block and prevent program or erase operations from changing the data in it. All blocks are locked after power-up or reset.

Two bus write cycles are required to issue the Block Lock command:

- The first bus cycle sets up the Block Lock command.
- The second bus write cycle latches the block address and locks the block.

The lock status can be monitored for each block using the Read Electronic Signature command. *Table 17* shows the lock status after issuing a Block Lock command.

Once set, the block lock bits remain set even after a hardware reset or power-down/powerup. They are cleared by a Block Unlock command.

Refer to *Section 9: Block locking* for a detailed explanation. See *Appendix C*, *Figure 30: Locking operations flowchart and pseudocode* for a flowchart for using the Lock command.

4.16 Block Unlock command

The Block Unlock command unlocks a block, allowing the block to be programmed or erased.

Two bus write cycles are required to issue the Block Unlock command:

- The first bus cycle sets up the Block Unlock command.
- The second bus write cycle latches the block address and unlocks the block.

The lock status can be monitored for each block using the Read Electronic Signature command. *Table 17* shows the protection status after issuing a Block Unlock command.

Refer to *Section 9: Block locking* for a detailed explanation and *Appendix C*, *Figure 30: Locking operations flowchart and pseudocode* for a flowchart for using the Block Unlock command.

4.17 Block Lock-Down command

The Block Lock-Down command is used to lock down a locked or unlocked block.

A locked-down block cannot be programmed or erased. The lock status of a locked-down block cannot be changed when WP is low, V_{IL} . When WP is high, V_{IH} , the lock-down function is disabled and the locked blocks can be individually unlocked by the Block Unlock command.

Two bus write cycles are required to issue the Block Lock-Down command:

- The first bus cycle sets up the Block Lock-Down command.
- The second bus write cycle latches the block address and locks down the block.

The lock status can be monitored for each block using the Read Electronic Signature command.

Locked-down blocks revert to the locked (and not locked-down) state when the device is reset on power-down. *Table 17* shows the lock status after issuing a Block Lock-Down command.

Refer to *Section 9: Block locking* for a detailed explanation and *Appendix C, Figure 30: Locking operations flowchart and pseudocode* for a flowchart for using the Lock-Down command.

N numonyx

		Bus operations							
Commands	Cycles		1st cycle		2nd cycle				
	S	Op.	Add	Data	Op.	Add	Data		
Read Array	1+	Write	BKA	FFh	Read	WA	RD		
Read Status Register	1+	Write	BKA	70h	Read	BKA ⁽²⁾	SRD		
Read Electronic Signature	1+	Write	BKA	90h	Read	BKA ⁽²⁾	ESD		
Read CFI Query	1+	Write	BKA	98h	Read	BKA ⁽²⁾	QD		
Clear Status Register	1	Write	Х	50h					
Block Erase	2	Write	BKA or BA ⁽³⁾	20h	Write	BA	D0h		
Program	2	Write	BKA or WA ⁽³⁾	40h or 10h	Write	WA	PD		
		Write	BA	E8h	Write	BA	n		
Buffer Program ⁽⁴⁾	n+4	Write	PA ₁	PD ₁	Write	PA ₂	PD ₂		
		Write	PA _{n+1}	PD _{n+1}	Write	Х	D0h		
Program/Erase Suspend	1	Write	Х	B0h					
Program/Erase Resume	1	Write	Х	D0h					
Protection Register Program	2	Write	PRA	C0h	Write	PRA	PRD		
Set Configuration Register	2	Write	CRD	60h	Write	CRD	03h		
Block Lock	2	Write	BKA or BA ⁽³⁾	60h	Write	BA	01h		
Block Unlock	2	Write	BKA or BA ⁽³⁾	60h	Write	BA	D0h		
Block Lock-Down	2	Write	BKA or BA ⁽³⁾	60h	Write	BA	2Fh		

Table 6.Standard commands⁽¹⁾

 X = Don't Care, WA = Word Address in targeted bank, RD =Read Data, SRD =Status Register Data, ESD = Electronic Signature Data, QD =Query Data, BA =Block Address, BKA = Bank Address, PD = Program Data, PRA = Protection Register Address, PRD = Protection Register Data, CRD = Configuration Register Data.

2. Must be same bank as in the first cycle. The signature addresses are listed in *Table 8*.

3. Any address within the bank can be used.

4. n+1 is the number of words to be programmed.

	Phase	Cycles	Bus write operations ⁽¹⁾									
Command			1st		2nd		3rd		Final -1		Final	
			Add	Data	Add	Data	Add	Data	Add	Data	Add	Data
Blank Check		2	BA	BCh	BA	CBh						
Buffer Enhanced Factory Program	Setup	2	BKA or WA ⁽²⁾	80h	WA ₁	D0h						
	Program/ verify ⁽³⁾	≥32	WA ₁	PD ₁	WA ₁	PD ₂	WA ₁	PD ₃	WA ₁	PD ₃₁	WA ₁	PD ₃₂
	Exit	1	NOT BA ₁ ⁽⁴⁾	х								

Table 7.Factory commands

1. WA = Word Address in targeted bank, BKA = Bank Address, PD =Program Data, BA = Block Address, X = 'don't care'.

2. Any address within the bank can be used.

3. The program/verify phase can be executed any number of times as long as the data is to be programmed to the same block.

4. WA₁ is the start address, NOT $BA_1 = Not Block Address of WA_1$.

Table 8.Electronic signature codes

	Code	Address (h)	Data (h)	
Manufacturer code		Bank address + 00	0020	
Device code	Тор	Bank address + 01	88C4 (M58LR128KT) 880D (M58LR256KT)	
Device code	Bottom	Bank address + 01	88C5 (M58LR128KB) 880E (M58LR256KB)	
	Locked		0001	
Block protection	Unlocked	Block address + 02	0000	
	Locked and locked-down	DIOCK AUDIESS + 02	0003	
	Unlocked and locked-down		0002	
Configuration Register		Bank address + 05	CR ⁽¹⁾	
Protection Register PR0 lock	Numonyx factory default		0002	
	OTP area permanently locked	Bank address + 80	0000	
Protection Projector PE	20	Bank address + 81 Bank address + 84	Unique device number	
Protection Register PF	IU	Bank address + 85 Bank address + 88	OTP Area	
Protection Register PF	R1 through PR16 Lock	Bank address + 89	PRLD ⁽¹⁾	
Protection Registers P	'R1-PR16	Bank address + 8A Bank address + 109	OTP area	

1. CR = Configuration Register, PRLD = Protection Register Lock Data.




		on nogiotoi				
	Lock		Description			
Number	Address Bits		Description			
	80h	Bit 0	Preprogrammed to protect unique device number, address 81h to 84h in PR0			
Lock 1		Bit 1	Protects 64 bits of OTP segment, address 85h to 88h in PR0			
		Bits 2 to 15	Reserved			
		Bit 0	Protects 128 bits of OTP segment PR1			
		Bit 1	Protects 128 bits of OTP segment PR2			
		Bit 2	Protects 128 bits of OTP segment PR3			
Lock 2	89h					
		Bit 13	Protects 128 bits of OTP segment PR14			
		Bit 14	Protects 128 bits of OTP segment PR15			
		Bit 15	Protects 128 bits of OTP segment PR16			

 Table 9.
 Protection Register locks

5 Status Register

The Status Register provides information on the current or previous program or erase operations. Issue a Read Status Register command to read the contents of the Status Register (refer to *Section 4.2: Read Status Register command* for more details). To output the contents, the Status Register is latched and updated on the falling edge of the Chip Enable or Output Enable signals and can be read until Chip Enable or Output Enable returns to V_{IH}. The Status Register can only be read using single Asynchronous or single synchronous reads. Bus read operations from any address within the bank always read the Status Register during program and erase operations if no Read Array command has been issued.

The various bits convey information about the status and any errors of the operation. Bits SR7, SR6, SR2 and SR0 give information on the status of the device and are set and reset by the device. Bits SR5, SR4, SR3 and SR1 give information on errors, they are set by the device but must be reset by issuing a Clear Status Register command or a hardware reset. If an error bit is set to '1' the Status Register should be reset before issuing another command.

The bits in the Status Register are summarized in *Table 10: Status Register bits*. Refer to *Table 10* in conjunction with the following text descriptions.

5.1 Program/Erase Controller status bit (SR7)

The Program/Erase Controller status bit indicates whether the Program/Erase Controller is active or inactive in any bank.

When the Program/Erase Controller status bit is Low (set to '0'), the Program/Erase Controller is active; when the bit is High (set to '1'), the Program/Erase Controller is inactive, and the device is ready to process a new command.

The Program/Erase Controller status bit is Low immediately after a Program/Erase Suspend command is issued until the Program/Erase Controller pauses. After the Program/Erase Controller pauses the bit is High.

5.2 Erase suspend status bit (SR6)

The erase suspend status bit indicates that an erase operation has been suspended in the addressed block. When the erase suspend status bit is High (set to '1'), a Program/Erase Suspend command has been issued and the memory is waiting for a Program/Erase Resume command.

The erase suspend status bit should only be considered valid when the Program/Erase Controller status bit is High (Program/Erase Controller inactive). SR6 is set within the erase suspend latency time of the Program/Erase Suspend command being issued, therefore, the memory may still complete the operation rather than entering the suspend mode.

When a Program/Erase Resume command is issued the erase suspend status bit returns Low.



5.3 Erase/blank check status bit (SR5)

The erase/blank check status bit identifies if there was an error during a block erase operation. When the erase/blank check status bit is High (set to '1'), the Program/Erase Controller has applied the maximum number of pulses to the block and still failed to verify that it has erased correctly.

The erase/blank check status bit should be read once the Program/Erase Controller status bit is High (Program/Erase Controller inactive).

The erase/blank check status bit is also used to indicate whether an error occurred during the blank check operation. If the data at one or more locations in the block where the Blank Check command has been issued is different from FFFFh, SR5 is set to '1'.

Once set High, the erase/blank check status bit must be set Low by a Clear Status Register command or a hardware reset before a new erase command is issued, otherwise the new command appears to fail.

5.4 **Program status bit (SR4)**

The program status bit identifies if there was an error during a program operation. It should be read once the Program/Erase Controller status bit is High (Program/Erase Controller inactive).

When the program status bit is High (set to '1'), the Program/Erase Controller has applied the maximum number of pulses to the word and still failed to verify that it has programmed correctly.

Attempting to program a '1' to an already programmed bit while $V_{PP} = V_{PPH}$ also sets the program status bit High. If V_{PP} is different from V_{PPH} , SR4 remains Low (set to '0') and the attempt is not shown.

Once set High, the program status bit must be set Low by a Clear Status Register command or a hardware reset before a new program command is issued, otherwise the new command appears to fail.

5.5 V_{PP} status bit (SR3)

The V_{PP} status bit identifies an invalid voltage on the V_{PP} pin during program and erase operations. The V_{PP} pin is only sampled at the beginning of a program or erase operation. Program and erase operations are not guaranteed if V_{PP} becomes invalid during an operation

When the V_{PP} status bit is Low (set to '0'), the voltage on the V_{PP} pin was sampled at a valid voltage.

When the V_{PP} status bit is High (set to '1'), the V_{PP} pin has a voltage that is below the V_{PP} lockout voltage, V_{PPLK}, the memory is protected and program and erase operations cannot be performed.

Once set High, the V_{PP} status bit must be set Low by a Clear Status Register command or a hardware reset before a new program or erase command is issued, otherwise the new command appears to fail.



5.6 **Program suspend status bit (SR2)**

The program suspend status bit indicates that a program operation has been suspended in the addressed block. The program suspend status bit should only be considered valid when the Program/Erase Controller status bit is High (Program/Erase Controller inactive).

When the program suspend status bit is High (set to '1'), a Program/Erase Suspend command has been issued and the memory is waiting for a Program/Erase Resume command.

SR2 is set within the program suspend latency time of the Program/Erase Suspend command being issued, therefore, the memory may still complete the operation rather than entering the suspend mode.

When a Program/Erase Resume command is issued the program suspend status bit returns Low.

5.7 Block protection status bit (SR1)

The block protection status bit identifies if a program or block erase operation has tried to modify the contents of a locked or locked-down block.

When the block protection status bit is High (set to '1'), a program or erase operation has been attempted on a locked or locked-down block

Once set High, the block protection status bit must be set Low by a Clear Status Register command or a hardware reset before a new program or erase command is issued, otherwise the new command appears to fail.

5.8 Bank write/multiple word program status bit (SR0)

The bank write status bit indicates whether the addressed bank is programming or erasing. In buffer enhanced factory program mode the multiple word program bit shows if the device is ready to accept a new word to be programmed to the memory array.

The bank write status bit should only be considered valid when the Program/Erase Controller status bit SR7 is Low (set to '0').

When both the Program/Erase Controller status bit and the bank write status bit are Low (set to '0'), the addressed bank is executing a program or erase operation. When the Program/Erase Controller status bit is Low (set to '0') and the bank write status bit is High (set to '1'), a program or erase operation is being executed in a bank other than the one being addressed.

In buffer enhanced factory program mode if multiple word program status bit is Low (set to '0'), the device is ready for the next word. If the multiple word program status bit is High (set to '1') the device is not ready for the next word.

For further details on how to use the Status Register, see the Flowcharts and Pseudocodes provided in *Appendix C*.



Bit	Name	Туре	Logic level ⁽¹⁾	Definition			
007	D/EC status	Chatwa	'1'	Ready			
SR7	P/EC status	Status	'0'	Busy			
SR6	Eraco succord status	Status	'1'	Erase sus	pended		
300	Erase suspend status	Sidius	'0'	Erase In p	rogress or completed		
SR5	Erase/blank check	Error	'1'	Erase/blank check error			
303	status	EIIU	'0'	Erase/blank check success Program error			
SR4	Program status	Error	'1'	Erase/blank check successProgram errorProgram success V_{PP} invalid, abort V_{PP} OKProgram suspendedProgram in progress or completedProgram/erase on protected block, abortNo operation to protected blocksSR7 = '1'Not allowedSR7 = '0'Program or erase operation in a bar other than the addressed bankSR7 = '1'No program or erase operation in the			
0114	T Togram status	LIIUI	'0'	BusyErase suspendedErase In progress or completedErase/blank check errorErase/blank check successProgram errorProgram success V_{PP} invalid, abort V_{PP} OKProgram in progress or completedProgram/erase on protected block, abortNo operation to protected blocksSR7 = '1'Not allowedSR7 = '1'SR7 = '0'Program or erase operation in the addressed bankSR7 = '1'Not allowedSR7 = '1'SR7 = '1'Not allowedSR7 = '1'Not allowedSR7 = '1'Not allowed			
SR3	V _{PP} status	Error	'1'	V _{PP} invalid, abort V _{PP} OK			
0110	Vpp Status	LIIOI	'0'	V _{PP} OK			
SR2	Program suspend	Status	'1'	Program s	uspended		
0112	status	Status	'0'	Program ir	Program in progress or completed		
SR1	Block protection	Error	'1'				
0111	status	LIIUI	'0'	No operation to protected blocks			
				SR7 = '1'	Not allowed		
			'1'	SR7 = '0'	Program or erase operation in a bank other than the addressed bank		
	Bank write status	Status	'0'	SR7 = '1'	No program or erase operation in the device		
			0	SR7 = '0'			
SR0				SR7 = '1'	Not allowed		
	Multiple word program status	Status	'1'	SR7 = '0'	The device is not ready for the next buffer loading or is going to exit the BEFP mode		
	(buffer enhanced factory program mode)		'0'	SR7 = '1' The device has exited the BEFP mode			
			U.	SR7 = '0'	The device is ready for the next buffer loading		

Table 10.Status Register bits

1. Logic level '1' is High, '0' is Low.

6 Configuration Register

The Configuration Register configures the type of bus access that the memory performs. Refer to *Section 7: Read modes* for details on read operations.

The Configuration Register is set through the command interface using the Set Configuration Register command. After a reset or power-up the device is configured for asynchronous read (CR15 = 1). The Configuration Register bits are described in *Table 12* They specify the selection of the burst length, burst type, burst X latency and the read operation. Refer to Figures 7 and 8 for examples of synchronous burst configurations.

6.1 Read select bit (CR15)

The read select bit, CR15, switches between asynchronous and synchronous read operations.

When the read select bit is set to '1', read operations are asynchronous, and when the read select bit is set to '0', read operations are synchronous.

Synchronous burst read is supported in both parameter and main blocks and can be performed across banks.

On reset or power-up the read select bit is set to '1' for asynchronous access.

6.2 X latency bits (CR13-CR11)

The X latency bits are used during synchronous read operations to set the number of clock cycles between the address being latched and the first data becoming available. Refer to *Figure 7: X latency and data output configuration example*.

For correct operation the X latency bits can only assume the values in *Table 12: Configuration Register.*

Table 11 shows how to set the X latency parameter, taking into account the speed class of the device and the frequency used to read the Flash memory in synchronous mode.

fmax	t _K min	X latency min
30 MHz	33 ns	3
40 MHz	25 ns	4
54 MHz	19 ns	5
66 MHz	15 ns	5

Table 11. X latency settings



6.3 Wait polarity bit (CR10)

The wait polarity bit is used to set the polarity of the Wait signal used in synchronous burst read mode. During synchronous burst read mode the Wait signal indicates whether the data output are valid or a WAIT state must be inserted.

When the wait polarity bit is set to '0' the Wait signal is active Low. When the wait polarity bit is set to '1' the Wait signal is active High.

6.4 Data output configuration bit (CR9)

The data output configuration bit configures the output to remain valid for either one or two clock cycles during synchronous mode.

When the data output configuration bit is '0' the output data is valid for one clock cycle, and when the data output configuration bit is '1' the output data is valid for two clock cycles.

The data output configuration bit must be configured using the following condition:

• $t_{K} > t_{KQV} + t_{QVK_CPU}$

where

- t_K is the clock period
- t_{QVK CPU} is the data setup time required by the system CPU
- t_{KQV} is the clock to data valid time.

If this condition is not satisfied, the data output configuration bit should be set to '1' (two clock cycles). Refer to *Figure 7: X latency and data output configuration example*.

6.5 Wait configuration bit (CR8)

The wait configuration bit is used to control the timing of the Wait output pin, WAIT, in synchronous burst read mode.

When WAIT is asserted, data is not valid and when WAIT is de-asserted, data is valid.

When the wait configuration bit is Low (set to '0') the Wait output pin is asserted during the WAIT state. When the wait configuration bit is High (set to '1'), the Wait output pin is asserted one data cycle before the WAIT state.

6.6 Burst type bit (CR7)

The burst type bit determines the sequence of addresses read during synchronous burst reads.

The burst type bit is High (set to '1'), as the memory outputs from sequential addresses only.

See *Table 13: Burst type definition* for the sequence of addresses output from a given starting address in sequential mode.

6.7 Valid Clock edge bit (CR6)

The valid Clock edge bit, CR6, configures the active edge of the Clock, K, during synchronous read operations. When the valid Clock edge bit is Low (set to '0') the falling edge of the Clock is the active edge. When the valid Clock edge bit is High (set to '1') the rising edge of the Clock is the active edge.

6.8 Wrap burst bit (CR3)

The wrap burst bit, CR3, selects between wrap and no wrap. Synchronous burst reads can be confined inside the 4, 8 or 16 word boundary (wrap) or overcome the boundary (no wrap).

When the Wrap Burst bit is Low (set to '0') the burst read wraps. When it is High (set to '1') the burst read does not wrap.

6.9 Burst length bits (CR2-CR0)

The burst length bits sets the number of words to be output during a synchronous burst read operation as result of a single address latch cycle.

They can be set for 4 words, 8 words, 16 words or continuous burst, where all the words are read sequentially. In continuous burst mode the burst sequence can cross bank boundaries.

In continuous burst mode, in 4, 8 or 16 words no-wrap, depending on the starting address, the device asserts the WAIT signal to indicate that a delay is necessary before the data is output.

If the starting address is shifted by 1, 2 or 3 positions from the four-word boundary, WAIT is asserted for 1, 2 or 3 clock cycles, respectively. When the burst sequence crosses the first 16-word boundary this indicates that the device needs an internal delay to read the successive words in the array. WAIT asserted only once during a continuous burst access. See also *Table 13: Burst type definition*.

CR14, CR5 and CR4 are reserved for future use.



Bit	Description	Value	Description
CR15	Read select	ValueDescription0Synchronous read1Asynchronous Read (default at power-on)1Asynchronous Read (default at power-on)0102 clock latency ⁽¹⁾ 0113 clock latency1004 clock latency1004 clock latency1015 clock latency1106 clock latency1117 clock latency (default)Other cligurations reserved0WAIT is active Low1Data held for one clock cycle1Data held for two clock cycles (default) ⁽¹⁾ 0WAIT is active during WAIT state1Sequential (default)0Reserved1Sequential (default)0Falling Clock edge1Rising Clock edge (default)0Wrap1No wrap (default)0014 words	
UNIS	neau seleci	1	Asynchronous Read (default at power-on)
CR14	Reserved		
		010	2 clock latency ⁽¹⁾
		011	3 clock latency
		100	4 clock latency
CR13-CR11	X latency	101	5 clock latency
		110	6 clock latency
		111	7 clock latency (default)
		Other co	onfigurations reserved
		0	WAIT is active Low
CR10	Wait polarity	1	WAIT is active High (default)
CD0		0 Data held for one clock cycle	
CR9	Data output configuration	1	Data held for two clock cycles (default) ⁽¹⁾
		0	WAIT is active during WAIT state
CR8	Wait configuration	1	WAIT is active one data cycle before WAIT state ⁽¹⁾ (default)
	Buret tures	0	Reserved
CR7	Burst type	1	Sequential (default)
CDC	Valid Clask adda	0	Falling Clock edge
CR6	Valid Clock edge	1	Rising Clock edge (default)
CR5-CR4	Reserved		
CR3	Wrap burst	0	Wrap
UNJ		1	No wrap (default)
		001	4 words
CR2-CR0	Burst Length	010	8 words
		011	16 words
		111	Continuous (default)

 Table 12.
 Configuration Register

1. The combination X latency=2, data held for two clock cycles and Wait active one data cycle before the WAIT state is not supported.

ιαυ	ie is.	Duistiyp	e definition				
Mode	Start		Sequer	ntial	Continuous burst		
Mo	addr	4 words	8 words	16 words			
	0	0-1-2-3	0-1-2-3-4-5-6- 7	0-1-2-3-4-5-6-7-8-9-10- 11-12-13-14-15	0-1-2-3-4-5-6		
	1	1-2-3-0	1-2-3-4-5-6-7- 0	1-2-3-4-5-6-7-8-9-10- 11-12-13-14-15-0	1-2-3-4-5-6-715-WAIT-16-17- 18		
	2	2-3-0-1	2-3-4-5-6-7-0- 1	2-3-4-5-6-7-8-9-10-11- 12-13-14-15-0-1	2-3-4-5-6-715-WAIT-WAIT- 16-17-18		
	3	3-0-1-2	3-4-5-6-7-0-1- 2	3-4-5-6-7-8-9-10-11-12- 13-14-15-0-1-2	3-4-5-6-715-WAIT-WAIT- WAIT-16-17-18		
Wrap							
N	7	7-4-5-6	7-0-1-2-3-4-5- 6	7-8-9-10-11-12-13-14- 15-0-1-2-3-4-5-6	7-8-9-10-11-12-13-14-15-WAIT- WAIT-WAIT-16-17		
	12				12-13-14-15-16-17-18		
	13				13-14-15-WAIT-16-17-18		
	14				14-15-WAIT-WAIT-16-17-18		
	15				15-WAIT-WAIT-WAIT-16-17- 18		

Table 13. Burst type definition

Mode	Start		Seque	ntial	Continuous burst
Mo	addr	4 words	8 words	16 words	Continuous burst
	0	0-1-2-3	0-1-2-3-4-5-6- 7	0-1-2-3-4-5-6-7-8-9-10- 11-12-13-14-15	
	1	1-2-3-4	1-2-3-4-5-6-7- 8	1-2-3-4-5-6-7-8-9-10- 11-12-13-14-15-WAIT- 16	
	2	2-3-4-5	2-3-4-5-6-7-8- 9	2-3-4-5-6-7-8-9-10-11- 12-13-14-15-WAIT- WAIT-16-17	
	3	3-4-5-6	3-4-5-6-7-8-9- 10	3-4-5-6-7-8-9-10-11-12- 13-14-15-WAIT-WAIT- WAIT-16-17-18	
No-wrap	7	7-8-9-10	7-8-9-10-11- 12-13-14	7-8-9-10-11-12-13-14- 15-WAIT-WAIT-WAIT- 16-17-18-19-20-21-22	Same as for wrap (wrap /no wrap has no effect on
No-				continuous burst)	
	12	12-13-14- 15	12-13-14-15- 16-17-18-19	12-13-14-15-16-17-18- 19-20-21-22-23-24-25- 26-27	
	13	13-14-15- WAIT-16	13-14-15- WAIT-16-17- 18-19-20	13-14-15-WAIT-16-17- 18-19-20-21-22-23-24- 25-26-27-28	
	14	14-15- WAIT- WAIT-16- 17	14-15-WAIT- WAIT-16-17- 18-19-20-21	14-15-WAIT-WAIT-16- 17-18-19-20-21-22-23- 24-25-26-27-28-29	
	15	15-WAIT- WAIT- WAIT-16- 17-18	15-WAIT- WAIT-WAIT- 16-17-18-19- 20-21-22	15-WAIT-WAIT- 16-17-18-19-20-21-22- 23-24-25-26-27-28-29- 30	

 Table 13.
 Burst type definition (continued)





Figure 7. X latency and data output configuration example

1. Amax is equal to A22 in the M58LR128KT/B and, to A23 in the M58LR256KT/B.

2. The settings shown are X latency = 4, data output held for one clock cycle.







1. Amax is equal to A22 in the M58LR128KT/B and, to A23 in the M58LR256KT/B.

7 Read modes

Read operations can be performed in two different ways depending on the settings in the Configuration Register. If the clock signal is 'don't care' for the data output, the read operation is asynchronous If the data output is synchronized with clock, the read operation is synchronous.

The read mode and format of the data output are determined by the Configuration Register. (see *Section 6: Configuration Register* for details). All banks support both asynchronous and synchronous read operations.

7.1 Asynchronous read mode

In asynchronous read operations the clock signal is 'don't care'. The device outputs the data corresponding to the address latched, that is the memory array, Status Register, common Flash interface or electronic signature depending on the command issued. CR15 in the Configuration Register must be set to '1' for asynchronous operations.

Asynchronous read operations can be performed in two different ways, Asynchronous random access read and asynchronous page read. Only asynchronous page read takes full advantage of the internal page storage so different timings are applied.

In asynchronous read mode a page of data is internally read and stored in a page buffer. The page has a size of 4 words and is addressed by address inputs A0 and A1.

The first read operation within the page has a longer access time (t_{AVQV} , random access time), subsequent reads within the same page have much shorter access times (t_{AVQV1} , page access time). If the page changes then the normal, longer timings apply again.

The device features an automatic standby mode. During asynchronous read operations, after a bus inactivity of 150 ns, the device automatically switches to automatic standby mode. In this condition the power consumption is reduced to the standby value and the outputs are still driven.

In asynchronous read mode, the WAIT signal is always deasserted.

See Table 24: Asynchronous read AC characteristics, Figure 11: Asynchronous random access read AC waveforms and Figure 12: Asynchronous page read AC waveforms for details.

7.2 Synchronous burst read mode

In synchronous burst read mode the data is output in bursts synchronized with the clock. It is possible to perform burst reads across bank boundaries.

Synchronous burst read mode can only be used to read the memory array. For other read operations, such as read Status Register, read CFI, and read electronic signature, single synchronous read or asynchronous random access read must be used.

In synchronous burst read mode the flow of the data output depends on parameters that are configured in the Configuration Register.

A burst sequence starts at the first clock edge (rising or falling depending on valid Clock edge bit CR6 in the Configuration Register) after the falling edge of Latch Enable or Chip Enable, whichever occurs last. Addresses are internally incremented and data is output on each data cycle after a delay which depends on the X latency bits CR13-CR11 of the Configuration Register.

The number of words to be output during a synchronous burst read operation can be configured as 4 words, 8 words, 16 words or continuous (burst length bits CR2-CR0). The data can be configured to remain valid for one or two clock cycles (data output configuration bit CR9).

The order of the data output can be modified through the wrap burst bit in the Configuration Register. The burst sequence is sequential and can be confined inside the 4, 8 or 16 word boundary (wrap) or overcome the boundary (no wrap).

The WAIT signal may be asserted to indicate to the system that an output delay will occur. This delay depends on the starting address of the burst sequence and on the burst configuration.

WAIT is asserted during the X latency, the WAIT state and at the end of a 4, 8 and 16 word burst. It is only de-asserted when output data are valid. In continuous burst read mode a WAIT state occurs when crossing the first 16 word boundary. If the starting address is aligned to the burst length (4, 8 or 16 words) the wrapped configuration has no impact on the output sequence.

The WAIT signal can be configured to be active Low or active High by setting CR10 in the Configuration Register.

See Table 25: Synchronous read AC characteristics and Figure 13: Synchronous burst read AC waveforms for details.



7.2.1 Synchronous burst read suspend

A synchronous burst read operation can be suspended, freeing the data bus for other higher priority devices. It can be suspended during the initial access latency time (before data is output) or after the device has output data. When the synchronous burst read operation is suspended, internal array sensing continues and any previously latched internal data is retained. A burst sequence can be suspended and resumed as often as required as long as the operating conditions of the device are met.

A synchronous burst read operation is suspended when Chip Enable, \overline{E} , is Low and the current address has been latched (on a Latch Enable rising edge or on a valid clock edge). The Clock signal is then halted at V_{II} or at V_{II}, and Output Enable, \overline{G} , goes High.

When Output Enable, \overline{G} , becomes Low again and the Clock signal restarts, the synchronous burst read operation is resumed exactly where it stopped.

WAIT reverts to high-impedance when Chip Enable, \overline{E} , or Output Enable, \overline{G} , goes High.

See Table 25: Synchronous read AC characteristics and Figure 15: Synchronous burst read suspend AC waveforms for details.

7.3 Single synchronous read mode

Single synchronous read operations are similar to synchronous burst read operations except that the memory outputs the same data to the end of the operation.

Synchronous single reads are used to read the electronic signature, Status Register, CFI, block protection status, Configuration Register Status or the Protection Register. When the addressed bank is in read CFI, read Status Register or read electronic signature mode, the WAIT signal is asserted during the X latency, the WAIT state and at the end of a 4, 8 and 16 word burst. It is only de-asserted when output data are valid.

See Table 25: Synchronous read AC characteristics and Figure 14: Single synchronous read AC waveforms for details.

8 Dual operations and multiple bank architecture

The multiple bank architecture of the M58LRxxxKT/B gives greater flexibility for software developers to split the code and data spaces within the memory array. The dual operations feature simplifies the software management of the device by allowing code to be executed from one bank while another bank is being programmed or erased.

The dual operations feature means that while programming or erasing in one bank, read operations are possible in another bank with zero latency (only one bank at a time is allowed to be in program or erase mode).

If a read operation is required in a bank, which is programming or erasing, the program or erase operation can be suspended.

In addition, if the suspended operation is erase then a program command can be issued to another block. This means the device can have one block in erase suspend mode, one programming, and other banks in read mode.

Bus read operations are allowed in another bank between setup and confirm cycles of program or erase operations.

By using a combination of these features, read operations are possible at any moment in the M58LRxxxKT/B device.

Dual operations between the parameter bank and either of the CFI, the OTP or the electronic signature memory space are not allowed. *Table 16* shows which dual operations are allowed or not between the CFI, the OTP, the electronic signature locations and the memory array.

Tables 14 and 15 show the dual operations possible in other banks and in the same bank.

			Con	nmands allo	wed in another bank				
Status of bank	Read Array	Read Status Register	Read CFI Query	Read Electronic Signature	Program, Buffer Program	Block Erase	Program /Erase Suspend	Program /Erase Resume	
Idle	Yes	Yes	Yes	Yes	Yes	Yes	Yes	Yes	
Programming	Yes	Yes	Yes	Yes	-	-	Yes	-	
Erasing	Yes	Yes	Yes	Yes	-	-	Yes	-	
Program suspended	Yes	Yes	Yes	Yes	-	-	-	Yes	
Erase suspended	Yes	Yes	Yes	Yes	Yes	_	_	Yes	

Table 14. Dual operations allowed in other banks



	Commands allowed in same bank									
Status of bank	Read Array	Read Status Register	Read CFI Query	Read Electronic Signature	Program, Buffer Program	Block Erase	Program /Erase Suspend	Program /Erase Resume		
Idle	Yes	Yes	Yes	Yes	Yes	Yes	Yes	Yes		
Programming	_(1)	Yes	Yes	Yes	-	-	Yes	-		
Erasing	_(1)	Yes	Yes	Yes	-	-	Yes	-		
Program suspended	Yes (2)	Yes	Yes	Yes	-	-	-	Yes		
Erase suspended	Yes ⁽²	Yes	Yes	Yes	Yes ⁽¹⁾	_	-	Yes		

 Table 15.
 Dual operations allowed in same bank

1. The Read Array command is accepted but the data output is not guaranteed until the program or erase has completed.

2. Not allowed in the block that is being erased or in the word that is being programmed.

Table 16. Dual operation limitations

	Programming/erasing parameter blocks bogramming/ asing main blocks Dotated in parameter bank Not located in parameter	Commands allowed					
Ourseast Status		Read CFI / OTP /	Read	Read ma	in blocks		
Current	Status	Electronic Signature	Parameter Blocks	Located in parameter bank	Not located in parameter bank		
•			No	No	Yes		
Programming/	parameter	Yes	No	No	Yes		
•		Yes	Yes	Yes	In different bank only		
Programming OTP		No	No	No	No		

9 Block locking

The M58LRxxxKT/B features an instant, individual block locking scheme that allows any block to be locked or unlocked with no latency. This locking scheme has three levels of protection.

- Lock/unlock this first level allows software only control of block locking.
- Lock-down this second level requires hardware interaction before locking can be changed.
- V_{PP} ≤ V_{PPLK} the third level offers complete hardware protection against program and erase on all blocks.

The protection status of each block can be set to locked, unlocked, and locked-down. *Table 17*, defines all of the possible protection states (\overline{WP} , DQ1, DQ0), and *Appendix C*, *Figure 30* shows a flowchart for the locking operations.

9.1 Reading block lock status

The lock status of every block can be read in read electronic signature mode of the device. To enter this mode issue the Read Electronic Signature command. Subsequent reads at the address specified in *Table 8* output the protection status of that block.

The lock status is represented by DQ0 and DQ1. DQ0 indicates the block lock/unlock status and is set by the Lock command and cleared by the Unlock command. DQ0 is automatically set when entering lock-down. DQ1 indicates the lock-down status and is set by the Lock-Down command. DQ1 cannot be cleared by software, except a hardware reset or power-down.

The following sections explain the operation of the locking system.

9.2 Locked state

The default status of all blocks on power-up or after a hardware reset is Locked (states (0,0,1) or (1,0,1)). Locked blocks are fully protected from program or erase operations. Any program or erase operations attempted on a locked block will return an error in the Status Register. The status of a locked block can be changed to unlocked or locked-down using the appropriate software commands. An unlocked block can be locked by issuing the Lock command.

9.3 Unlocked state

Unlocked blocks (states (0,0,0), (1,0,0) (1,1,0)), can be programmed or erased. All unlocked blocks return to the locked state after a hardware reset or when the device is powered-down. The status of an unlocked block can be changed to locked or locked-down using the appropriate software commands. A locked block can be unlocked by issuing the Unlock command.

9.4 Lock-down state

Blocks that are locked-down (state (0,1,x))are protected from program and erase operations (as for locked blocks) but their protection status cannot be changed using software commands alone. A locked or unlocked block can be locked down by issuing the Lock-Down command. Locked-down blocks revert to the locked state when the device is reset or powered-down.

The lock-down function is dependent on the Write Protect, WP, input pin.

When $\overline{WP}=0$ (V_{IL}), the blocks in the lock-down state (0,1,x) are protected from program, erase and protection status changes.

When $\overline{WP}=1$ (V_{IH}) the lock-down function is disabled (1,1,x) and locked-down blocks can be individually unlocked to the (1,1,0) state by issuing the software command, where they can be erased and programmed.

When the lock-down function is disabled (\overline{WP} =1) blocks can be locked (1,1,1) and unlocked (1,1,0) as desired. When \overline{WP} =0 blocks that were previously locked-down return to the lock-down state (0,1,x) regardless of any changes that were made while \overline{WP} =1.

Device reset or power-down resets all blocks, including those in lock-down, to the locked state.

9.5 Locking operations during erase suspend

Changes to block lock status can be performed during an erase suspend by using the standard locking command sequences to unlock, lock or lock-down a block. This is useful in the case when another block needs to be updated while an erase operation is in progress.

To change block locking during an erase operation, first write the Erase Suspend command, then check the Status Register until it indicates that the erase operation has been suspended. Next write the desired Lock command sequence to a block and the lock status is changed. After completing any desired lock, read, or program operations, resume the erase operation with the Erase Resume command.

If a block is locked or locked down during an erase suspend of the same block, the locking status bits is changed immediately, but when the erase is resumed, the erase operation completes. Locking operations cannot be performed during a program suspend.

-	tection status ⁽¹⁾ DQ1, DQ0)	Next protection status ⁽¹⁾ (WP, DQ1, DQ0)						
Current Program/erase allowed		After Block Lock command	After Block Unlock command	After Block Lock-Down command	After WP transition			
1,0,0	yes	1,0,1	1,0,0	1,1,1	0,0,0			
1,0,1 ⁽²⁾	no	1,0,1	1,0,0	1,1,1	0,0,1			
1,1,0	yes	1,1,1	1,1,0	1,1,1	0,1,1			
1,1,1	no	1,1,1	1,1,0	1,1,1	0,1,1			
0,0,0	yes	0,0,1	0,0,0	0,1,1	1,0,0			
0,0,1 ⁽²⁾	no	0,0,1	0,0,0	0,1,1	1,0,1			
0,1,1	no	0,1,1	0,1,1	0,1,1	1,1,1 or 1,1,0 ⁽³⁾			

Table 17.Lock status

1. The lock status is defined by the write protect pin and by DQ1 ('1' for a locked-down block) and DQ0 ('1' for a locked block) as read in the Read Electronic Signature command with DQ1 = V_{IH} and DQ0 = V_{IL} .

2. All blocks are locked at power-up, so the default configuration is 001 or 101 according to $\overline{\text{WP}}$ status.

3. A $\overline{\text{WP}}$ transition to V_{IH} on a locked block will restore the previous DQ0 value, giving a 111 or 110.



10 Program and erase times and endurance cycles

The program and erase times and the number of program/ erase cycles per block are shown in *Table 18.* Exact erase times may change depending on the memory array condition. The best case is when all the bits in the block are at '0' (pre-programmed). The worst case is when all the bits in the block are at '1' (not preprogrammed). Usually, the system overhead is negligible with respect to the erase time. In the M58LRxxxKT/B the maximum number of program/erase cycles depends on the V_{PP} voltage supply used.

	Parameter	Co	ondition	Min	Тур	Typical after 100kW/E cycles	Max	Unit
		Parameter block	(16 Kword)		0.6	1	2.5	S
	Erase	Main block (64	Preprogrammed		1.2	3	4	S
		Kword)	Not preprogrammed		1.5		4	s
	D (3)	Single Word	Word program		12		180	μs
V _{DD}			Buffer program		12		180	μs
Ш	Program ⁽³⁾	Buffer (32 words) (buffer program)			384			μs
V_{PP}		Main block (64 Kword)			768			ms
	Suspend latency	Program	Program		20		25	μs
	Suspend latency	Erase			20		25	μs
	Program/erase cycles	Main blocks	Main blocks					cycles
	(per block)	Parameter blocks		100 000				cycles

 Table 18.
 Program/erase times and endurance cycles^{(1) (2)}

	· · · · · ·	e Parameter block (16 Kword) Main block (64 Kword) Single word Buffer enhanced factory program Buffer (32 words) Buffer enhanced factory program		1	,			
	Parameter	Coi	ndition	Min	Тур	Typical after 100kW/E cycles	Max	Unit
	Erase	Parameter block (16 Kword)		0.6		2.5	s
	Elase	Main block (64 Kw	vord)		1		4	s
		Single word	Word program		10		170	μs
			Buffer enhanced factory program ⁽⁴⁾		2.5			μs
		Buffer (32 words)	Buffer program		80			μs
Т	Program ⁽³⁾				80			μs
>		Main block (64 Kwords)	Buffer program		160			ms
V _{PP} =			Buffer enhanced factory program		160			ms
			Buffer program		1.28			s
		Bank (8 Mbits)	Buffer enhanced factory program		1.28			s
	Program/erase cycles	Main blocks	Main blocks				1000	cycles
	(per block)	Parameter blocks				2500	cycles	
	Blank Chaok	Main blocks			16			ms
	Blank Check	Parameter blocks			4			ms

Table 18.	Program/erase times and endurance cycles ^{(1) (2)} (continued
Table To.	Frogram/erase times and endurance cycles, / / (continue

1. $T_A = -30$ to 85 °C; $V_{DD} = 1.7$ V to 2 V; $V_{DDQ} = 1.7$ V to 2 V.

2. Values are liable to change with the external system-level overhead (command sequence and Status Register polling execution).

3. Excludes the time needed to execute the command sequence.

4. This is an average value on the entire device.

11 Maximum ratings

Stressing the device above the rating listed in the *Table 19: Absolute maximum ratings* may cause permanent damage to the device. These are stress ratings only and operation of the device at these or any other conditions above those indicated in the operating sections of this specification is not implied. Exposure to Absolute maximum rating conditions for extended periods may affect device reliability. Refer also to the Numonyx SURE program and other relevant quality documents.

Symbol	Parameter	Va	lue	Unit
Symbol	Falameter	Min	Мах	Unit
Τ _Α	Ambient operating temperature	-30	85	°C
T _{BIAS}	Temperature under bias	-30	85	°C
T _{STG}	Storage temperature	-65	125	°C
V _{IO}	Input or output voltage	-0.5	V _{DDQ} + 0.6	V
V _{DD}	Supply voltage	-0.2	2.5	V
V _{DDQ}	Input/output supply voltage	-0.2	2.5	V
V _{PP}	Program voltage	-0.2	10	V
Ι _Ο	Output short circuit current		100	mA
t _{VPPH}	Time for V_{PP} at V_{PPH}		100	hours

 Table 19.
 Absolute maximum ratings

12 DC and AC parameters

This section summarizes the operating measurement conditions and the DC and AC characteristics of the device. The parameters in the DC and AC characteristics tables in this section are derived from tests performed under the measurement conditions summarized in *Table 20: Operating and AC measurement conditions*. Designers should check that the operating conditions in their circuit match the operating conditions when relying on the quoted parameters.

Parameter	70) ns	85	Unit	
	Min	Max	Min	Max	
V _{DD} supply voltage	1.7	2.0	1.7	2.0	V
V _{DDQ} supply voltage	1.7	2.0	1.7	2.0	V
V _{PP} supply voltage (factory environment)	8.5	9.5	8.5	9.5	V
V _{PP} supply voltage (application environment)	-0.4	V _{DDQ} +0.4	-0.4	V _{DDQ} +0.4	V
Ambient operating temperature	-30	85	-30	85	°C
Load capacitance (CL)	30		:	30	pF
Input rise and fall times		5		5	ns
Input pulse voltages	0 to V _{DDQ}		0 to	V	
Input and output timing ref. voltages	VD	_{DQ} /2	VD	_{DQ} /2	V

Table 20.	Operating and AC measurement conditions
-----------	---

Figure 9. AC measurement I/O waveform







Figure 10. AC measurement load circuit

Table 21. Capacitance⁽¹⁾

Symbol	Parameter	Test condition	Min	Max	Unit
C _{IN}	Input capacitance	V _{IN} = 0 V	6	8	pF
C _{OUT}	Output capacitance	V _{OUT} = 0 V	8	12	pF

1. Sampled only, not 100% tested.

Symbol	Paramete	r	Test condition	Тур	Max	Unit
I _{LI}	Input leakage current		$0~V \leq V_{IN} \leq V_{DDQ}$		±1	μA
I _{LO}	Output leakage current		$0 V \le V_{OUT} \le V_{DDQ}$		±1	μΑ
	Supply current Asynchronous read (f = 5 MHz)		$\overline{E} = V_{IL}, \overline{G} = V_{IH}$	13	15	mA
			4 word	18	20	mA
	Supply current		8 word	20	22	mA
	Synchronous read (f = 5	54 MHz)	16 word	25	27	mA
I _{DD1}			Continuous	28	30	mA
			4 word	20	22	mA
	Supply current		8 word	22	24	mA
	Synchronous read (f = 6	6 MHz)	16 word	27	29	mA
			Continuous	30	32	mA
		128 Mbit		22	70	μA
I _{DD2}	Supply current (reset)	256 Mbit	$\overline{RP} = V_{SS} \pm 0.2 V$	70	110	
	Supply current	128 Mbit	$\overline{E} = V_{DD} \pm 0.2 V$	22	70	
I _{DD3}	(standby)	256 Mbit	K=V _{SS}	70	110	μA
I _{DD4}	Supply current (automa	tic standby)	$\overline{E} = V_{IL}, \overline{G} = V_{IH}$	22	50	μΑ
	Supply ourrept (program	2)	V _{PP} = V _{PPH}	10	30	mA
ı (1)	Supply current (program	1)	$V_{PP} = V_{DD}$	20	35	mA
$I_{DD5}^{(1)}$	Supply ourrant (araga)		V _{PP} = V _{PPH}	10	30	mA
	Supply current (erase)		$V_{PP} = V_{DD}$	20	35	mA
	Supply ourrent		Program/erase in one bank, asynchronous read in another bank	33	50	mA
I _{DD6} ^{(1),(2)}	Supply current (dual operations)		Program/erase in one bank, synchronous read (continuous f = 66 MHz) in another bank	50	67	mA
. (1)	Supply current	128 Mbit	$\overline{E} = V_{DD} \pm 0.2 V$	22	70	
I _{DD7} ⁽¹⁾	program/erase suspended (standby)	256 Mbit	K=V _{SS}	70	110	μA
	V _{PP} supply current (pro	oram)	V _{PP} = V _{PPH}	2	5	mA
I _{PP1} ⁽¹⁾		granny	$V_{PP} = V_{DD}$	0.2	5	μΑ
'PP1` ′	V _{PP} supply current (eras	<u> </u>	V _{PP} = V _{PPH}	2	5	mA
		30)	$V_{PP} = V_{DD}$	0.2	5	μA

Table 22. DC characteristics - currents



Symbol	Parameter		Test condition	Тур	Max	Unit
I _{PP2}	V _{PP} supply current (read)		$V_{PP} \leq V_{DD}$	0.2	5	μΑ
I _{PP3} ⁽¹⁾	V _{PP} supply current (standby)		$V_{PP} \le V_{DD}$	0.2	5	μΑ

Table 22. DC characteristics - currents

1. Sampled only, not 100% tested.

2. V_{DD} dual operation current is the sum of read and program or erase currents.

Table 23. DC characteristics - voltages

Symbol	Parameter	Test condition	Min	Тур	Мах	Unit
V _{IL}	Input Low voltage		0		0.4	V
V _{IH}	Input High voltage		V _{DDQ} -0.4		V _{DDQ} + 0.4	V
V _{OL}	Output Low voltage	I _{OL} = 100 μA			0.1	V
V _{OH}	Output High voltage	I _{OH} = −100 μA	V _{DDQ} –0.1			V
V _{PP1}	V _{PP} Program voltage-logic	Program, erase	1.3	1.8	3.3	V
V _{PPH}	V _{PP} Program voltage factory	Program, erase	8.5	9.0	9.5	V
V _{PPLK}	Program or erase lockout				0.4	V
V _{LKO}	V _{DD} lock voltage				1	V





Figure 11. Asynchronous random access read AC waveforms





Figure 12. Asynchronous page read AC waveforms

	um h a l	A 14	Devenueter		M58LR>	M58LRxxxKT/B	
Symbol		Alt	Parameter	70	85	Unit	
	t _{AVAV}	t _{RC}	Address Valid to Next Address Valid	Min	70	85	ns
	t _{AVQV}	t _{ACC}	Address Valid to Output Valid (Random)	Max	70	85	ns
	t _{AVQV1}	t _{PAGE}	Address Valid to Output Valid (page)	Max	20	25	ns
	$t_{AXQX}^{(1)}$	t _{OH}	Address Transition to Output Transition	Min	0	0	ns
	t _{ELTV}		Chip Enable Low to Wait Valid	Max	11	14	ns
	$t_{ELQV}^{(2)}$	t _{CE}	Chip Enable Low to Output Valid	Max	70	85	ns
sbi	$t_{\text{ELQX}}^{(1)}$	t _{LZ}	Chip Enable Low to Output Transition	Min	0	0	ns
Read Timings	t _{EHTZ}		Chip Enable High to Wait Hi-Z	Max	11	14	ns
ad T	$t_{\text{EHQX}}^{}^{(1)}$	t _{OH}	Chip Enable High to Output Transition	Min	2	2	ns
Re	$t_{\text{EHQZ}}^{(1)}$	t _{HZ}	Chip Enable High to Output Hi-Z	Max	14	14	ns
	$t_{GLQV}^{(2)}$	t _{OE}	Output Enable Low to Output Valid	Max	20	20	ns
	$t_{\text{GLQX}}^{\left(1\right)}$	t _{OLZ}	Output Enable Low to Output Transition	Min	0	0	ns
	t _{GLTV}		Output Enable Low to Wait Valid	Max	14	14	ns
	${\rm t}_{\rm GHQX}^{(1)}$	t _{OH}	Output Enable High to Output Transition	Min	2	2	ns
	${\rm t_{GHQZ}}^{(1)}$	t _{DF}	Output Enable High to Output Hi-Z	Max	14	14	ns
	t _{GHTZ}		Output Enable High to Wait Hi-Z	Max	14	14	ns
	t _{AVLH}	t _{AVADVH}	Address Valid to Latch Enable High	Min	7	7	ns
sɓu	t _{ELLH}	t _{ELADVH}	Chip Enable Low to Latch Enable High	Min	10	10	ns
Timii	t _{LHAX}	t _{ADVHAX}	Latch Enable High to Address Transition	Min	7	7	ns
Latch Timings	t _{LLLH}	t _{ADVLADVH}	Latch Enable Pulse Width	Min	7	7	ns
La	t _{LLQV}	t _{ADVLQV}	Latch Enable Low to Output Valid (Random)	Max	70	85	ns

 Table 24.
 Asynchronous read AC characteristics

1. Sampled only, not 100% tested.

2. \overline{G} may be delayed by up to t_{ELQV} - t_{GLQV} after the falling edge of \overline{E} without increasing t_{ELQV} .

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Figure 13. Synchronous burst read AC waveforms



Figure 14. Single synchronous read AC waveforms

1. Amax is equal to A22 in the M58LR128KT/B and, to A23 in the M58LR256KT/B.

2. Address latched and data output on the rising clock edge. Either the rising or the falling edge of the clock signal, K, can be configured as the active edge. Here, the active edge is the rising one.

3. The WAIT signal is configured to be active during wait state. WAIT signal is active Low.





Figure 15. Synchronous burst read suspend AC waveforms



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Table 25.	Synchronous read AC characteristics ^{(1) (2)}
-----------	--

Symbol		Alt	Parameter		M58LR>	xxKT/B	Unit
		Alt	Parameter	Parameter			Unit
	t _{AVKH}	t _{AVCLKH}	Address Valid to Clock High	Min	7	7	ns
	t _{ELKH}	t _{ELCLKH}	Chip Enable Low to Clock High	Min	7	7	ns
gs	t _{ELTV}		Chip Enable Low to Wait Valid	Max	14	14	ns
Read Timings	t _{EHEL}		Chip Enable Pulse Width (subsequent synchronous reads)	Min	14	14	ns
Rea	t _{EHTZ}		Chip Enable High to Wait Hi-Z	Max	14	14	ns
sno	t _{KHAX}	t _{CLKHAX}	Clock High to Address Transition	Min	7	7	ns
Synchronous	t _{KHQV} t _{KHTV}	t _{CLKHQV}	Clock High to Output Valid Clock High to WAIT Valid	Max	11	14	ns
S	t _{KHQX} t _{KHTX}	t _{CLKHQX}	Clock High to Output Transition Clock High to WAIT Transition	Min	3	3	ns
	t _{LLKH}	t _{ADVLCLKH}	Latch Enable Low to Clock High	Min	7	7	ns
su	+	+	Clock Period (f = 54 MHz)	Min		18.5	ns
catio	^t кнкн	t _{CLK}	Clock Period (f = 66 MHz)	IVIIII	15		ns
Specifications	t _{KHKL} t _{KLKH}		Clock High to Clock Low Clock Low to Clock High	Min	4.5	4.5	ns
Clock	t _f t _r		Clock Fall or Rise Time	Max	3	3	ns

1. Sampled only, not 100% tested.

2. For other timings please refer to Table 24: Asynchronous read AC characteristics.


Figure 17. Write AC waveforms, Write Enable controlled

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		Alt		M58LRxxxKT/B				
S	Symbol		Parameter			85	Unit	
	t _{AVAV}	t _{WC}	Address Valid to Next Address Valid	Min	70	85	ns	
	t _{AVLH}		Address Valid to Latch Enable High	Min	7	7	ns	
	t _{AVWH} ⁽²⁾		ddress Valid to Write Enable High Min		45	45	ns	
	t _{DVWH}	t _{DS}	Data Valid to Write Enable High	Min	45	45	ns	
	t _{ELLH}		Chip Enable Low to Latch Enable High	Min	10	10	ns	
	t _{ELWL}	t _{CS}	Chip Enable Low to Write Enable Low	Min	0	0	ns	
sɓu	t _{ELQV}		Chip Enable Low to Output Valid	Min	70	85	ns	
Limir	t _{ELKV}		Chip Enable Low to Clock Valid	Min	7	7	ns	
led ⁻	t _{GHWL}		Output Enable High to Write Enable Low	Min	17	17	ns	
ntrol	t _{LHAX}		Latch Enable High to Address Transition	Min	7	7	ns	
°C €	t _{LLLH}		Latch Enable Pulse Width Min		7	7	ns	
abl€	t _{WHAV} ⁽²⁾	(2) Write Enable High to Address Valid Min		0	0	ns		
Write Enable Controlled Timings	t _{WHAX} ⁽²⁾	t _{AH}	Write Enable High to Address Transition	Min	0	0	ns	
Writ	t _{WHDX}	t _{DH}	Write Enable High to Input Transition	Min	0	0	ns	
	t _{WHEH}	H t _{CH} Write Enable High to Chip Enable High		Min	0	0	ns	
	t _{WHEL} ⁽³⁾		Write Enable High to Chip Enable Low	Min	25	25	ns	
	t _{WHGL}		Write Enable High to Output Enable Low	Min	0	0	ns	
	t _{WHLL} ⁽³⁾		Write Enable High to Latch Enable Low	Min	25	25	ns	
	t _{WHWL}	t _{WPH}	Write Enable High to Write Enable Low	Min	25	25	ns	
	t _{WLWH}	t _{WP}	Write Enable Low to Write Enable High	Min	45	45	ns	
	t _{QVVPL}		Output (Status Register) Valid to V_{PP} Low	Min	0	0	ns	
Protection Timings	t _{QVWPL}		Output (Status Register) Valid to Write Protect Low	Min	0	0	ns	
Du T	t _{VPHWH}	t _{VPS}	V _{PP} High to Write Enable High	Min	200	200	ns	
tectic	t _{WHVPL}		Write Enable High to V _{PP} Low	Min	200	200	ns	
Proi	t _{WHWPL}		Write Enable High to Write Protect Low	Min	200	200	ns	
	t _{WPHWH}		Write Protect High to Write Enable High	Min	200	200	ns	

 Table 26.
 Write AC characteristics, Write Enable controlled⁽¹⁾

1. Sampled only, not 100% tested.

2. Meaningful only if \overline{L} is always kept low.

3. t_{WHEL} and t_{WHLL}have this value when reading in the targeted bank or when reading following a Set Configuration Register command. System designers should take this into account and may insert a software No-Op instruction to delay the first read in the same bank after issuing any command and to delay the first read to any address after issuing a Set Configuration Register command. If the first read after the command is a Read Array operation in a different bank and no changes to the Configuration Register have been issued, t_{WHEL} and t_{WHLL} is 0 ns.



Figure 18. Write AC waveforms, Chip Enable controlled

Symbol Alt		hol Alt Beromotor			M58LR>	xxKT/B	11
		Alt	Parameter		70	85	Unit
	t _{AVAV}	t _{WC}	Address Valid to Next Address Valid	Min	70	85	ns
	t _{AVEH}		Address Valid to Chip Enable High	Min	45	45	ns
	t _{AVLH}		ddress Valid to Latch Enable High Min		7	7	ns
	t _{DVEH}	t _{DS}	ata Valid to Chip Enable High Min		45	45	ns
	t _{EHAX}	t _{AH}	Chip Enable High to Address Transition	Min	0	0	ns
sối	t _{EHDX}	t _{DH}	Chip Enable High to Input Transition	Min	0	0	ns
Chip Enable Controlled Timings	t _{EHEL}	t _{CPH}	Chip Enable High to Chip Enable Low	Min	25	25	ns
led T	t _{EHGL}		Chip Enable High to Output Enable Low	Min	0	0	ns
ntrol	t _{EHWH}	t _{CH}	Chip Enable High to Write Enable High	Min	0	0	ns
CO	t _{ELKV}		Chip Enable Low to Clock Valid Min		7	7	ns
lable	t _{ELEH}	t _{CP}	Chip Enable Low to Chip Enable High Min		45	45	ns
рĒ	t _{ELLH}		Chip Enable Low to Latch Enable High Min		10	10	ns
Chi	t _{ELQV}		Chip Enable Low to Output Valid Min		70	85	ns
	t _{GHEL}		Output Enable High to Chip Enable Low Min		17	17	ns
	t _{LHAX}		Latch Enable High to Address Transition Min		7	7	ns
	t _{LLLH}		Latch Enable Pulse Width Min		7	7	ns
	t _{WHEL} ⁽²⁾		Write Enable High to Chip Enable Low	Min	25	25	ns
	t _{WLEL}	t _{CS}	Write Enable Low to Chip Enable Low	Min	0	0	ns
	t _{EHVPL}		Chip Enable High to V _{PP} Low	Min	200	200	ns
sɓu	t _{EHWPL}		Chip Enable High to Write Protect Low	Min	200	200	ns
Timi	t _{QVVPL}		utput (Status Register) Valid to V _{PP} Low Min		0	0	ns
Protection Timings	t _{QVWPL}		Output (Status Register) Valid to Write Protect Low Min		0	0	ns
Proi	t _{VPHEH}	t _{VPS}	V _{PP} High to Chip Enable High	Min	200	200	ns
	t _{WPHEH}		Write Protect High to Chip Enable High	Min	200	200	ns

Table 27. W	Irite AC characteristics, Chip Enable controlled ⁽¹⁾
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1. Sampled only, not 100% tested.

2. t_{WHEL} has this value when reading in the targeted bank or when reading following a Set Configuration Register command. System designers should take this into account and may insert a software No-Op instruction to delay the first read in the same bank after issuing any command and to delay the first read to any address after issuing a Set Configuration Register command. If the first read after the command is a Read Array operation in a different bank and no changes to the Configuration Register have been issued, t_{WHEL} is Ons.



Figure 19. Reset and power-up AC waveforms

Table 28.	Reset and I	power-up A	AC characteristics
	neoci una j		

Symbol	Parameter Test condition		M58LRxxxKT/B	Unit	
t _{PLWL}	Reset Low to Write Enable Low	During program	Min	25	μs
t _{PLEL}	Reset Low to Chip Enable Low,	During erase	Min	25	μs
t _{PLGL} t _{PLLL}	Reset Low to Output Enable Low, Reset Low to Latch Enable Low	Other conditions	Min	80	ns
t _{PHWL} t _{PHEL} t _{PHGL} t _{PHLL}	Reset High to Write Enable Low Reset High to Chip Enable Low Reset High to Output Enable Low Reset High to Latch Enable Low		Min	30	ns
t _{PLPH} ^{(1),(2)}	RP pulse width		Min	50	ns
t _{VDHPH} ⁽³⁾	Supply voltages High to Reset High		Min	300	μs

1. The device Reset is possible but not guaranteed if t_{PLPH} < 50 ns.

2. Sampled only, not 100% tested.

3. It is important to assert \overline{RP} to allow proper CPU initialization during power-up or Reset.

13 Package mechanical

To meet environmental requirements, Numonyx offers the M58LR128KT/B in ECOPACK® packages. These packages have a lead-free, second-level interconnect. The category of second-level interconnect is marked on the package and on the inner box label, in compliance with JEDEC Standard JESD97.

The maximum ratings related to soldering conditions are also marked on the inner box label. ECOPACK is an Numonyx trademark.

Figure 20. VFBGA56 7.7 x 9 mm - 8 x 7 active ball array, 0.75 mm pitch, bottom view package outline



1. Drawing is not to scale.



	mechanical	data				
Cumhal		Millimeters				
Symbol	Тур	Min	Max	Тур	Min	Max
А			1.000			0.0394
A1		0.200			0.0079	
A2	0.660			0.0260		
b	0.350	0.300	0.400	0.0138	0.0118	0.0157
D	7.700	7.600	7.800	0.3031	0.2992	0.3071
D1	5.250	-	-	0.2067	-	-
ddd			0.080			0.0031
е	0.750	-	-	0.0295	-	-
E	9.000	8.900	9.100	0.3543	0.3504	0.3583
E1	4.500	-	-	0.1772	-	-
FD	1.225	-	-	0.0482	-	-
FE	2.250	-	_	0.0886	-	-
SD	0.375	-	-	0.0148	-	-

 Table 29.
 VFBGA56 7.7 x 9 mm - 8 x 7 ball array, 0.75 mm pitch, package mechanical data

Figure 21. TFBGA88 8 x 10 mm - 8 x 10 ball array, 0.8 mm pitch, bottom view package outline



Figure 22. Drawing is not to scale.



	uala					
Symbol		Millimeters		Inches		
Symbol	Тур	Min	Max	Тур	Min	Max
А			1.200			0.0472
A1		0.200			0.0079	
A2	0.850			0.0335		
b	0.350	0.300	0.400	0.0138	0.0118	0.0157
D	8.000	7.900	8.100	0.3150	0.3110	0.3189
D1	5.600			0.2205		
ddd			0.100			0.0039
E	10.000	9.900	10.100	0.3937	0.3898	0.3976
E1	7.200			0.2835		
E2	8.800			0.3465		
е	0.800	-	-	0.0315	-	-
FD	1.200			0.0472		
FE	1.400			0.0551		
FE1	0.600			0.0236		
SD	0.400			0.0157		
SE	0.400			0.0157		

Table 30.TFBGA88 8 x 10 mm - 8 x 10 ball array, 0.8 mm pitch, package mechanical
data

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Figure 23. VFBGA79 9x11mm - 0.75 pitch -package outline ⁽¹⁾

1. Drawing not to scale

		11111 - 0.75 11111	pitcii, packag		uata		
Symbol		Millimeters		Inches			
Symbol	Min	Тур	Max	Min	Тур	Мах	
А			1.000			0.0394	
A1	0.150			0.0059			
A2		0.665			0.0262		
b	0.325	0.375	0.425	0.0128	0.0148	0.0167	
D	10.900	11.000	11.100	0.4291	0.4331	0.4370	
Е	8.900	9.000	9.100	0.3504	0.3543	0.3583	
е		0.750			0.0295		
Ν		79			79		
Y			0.100			0.0039	
S1	0.900	1.000	1.100	0.0354	0.0394	0.0433	
S2	2.150	2.250	2.350	0.0846	0.0886	0.0925	

Table 31. VFBGA79 9x11mm - 0.75 mm pitch, package mechanical data

14 Part numbering

Table 32. Ordering information scheme

Example:	M58LR128KT	70 ZB 5 E
Device type		
M58		
Architecture		
L = multilevel, multiple bank, burst mode		
Operating voltage		
$R = V_{DD} = 1.7 \text{ V to } 2.0 \text{ V}, V_{DDQ} = 1.7 \text{ V to } 2.0 \text{ V}$		
Density		
128 = 128 Mbit (x16)		
256 = 256 Mbit (x16)		
Technology		
K = 65 nm technology		
Parameter location		
T = top boot		
B = bottom boot		
Speed		
70 = 70 ns		
85 = 85 ns		
- .		
ZB = VFBGA56, 7.7 x 9 mm, 0.75 mm pitch ⁽¹⁾		
ZC = VFBGA79, 9 x 11 mm, 0.75 mm pitch		
ZQ = TFBGA88, 8 x 10mm, 0.8 mm pitch		
Temperature range		
5 = -30 to 85 °C		
Packing option		

E = ECOPACK® package, standard packing

F = ECOPACK® package, tape and reel packing

1. The M58LR128KT/B are available in the ZB package, while the M58LR256KT/B are available M58LR256KT/B are available in the ZC and ZQ package.

Devices are shipped from the factory with the memory content bits erased to '1'.

For a list of available options (speed, package, etc.) or for further information on any aspect of this device, please contact the Numonyx sales office nearest to you.

Appendix A Block address tables

The following set of equations can be used to calculate a complete set of block addresses for the M58LRxxxKT/B using the information contained in Tables *36* to *44*.

To calculate the block base address from the block number:

First it is necessary to calculate the bank number and the block number offset. This can be achieved using the following formulas:

 $Bank_Number = (Block_Number - 3) / 8$

Block_Number_Offset = Block_Number - 3 - (Bank_Number x 8),

If Bank_Number= 0, the block base address can be directly read from Tables *36* and *42* (parameter bank block addresses) in the address range column, in the row that corresponds to the given block number.

Otherwise:

Block_Base_Address = Bank_Base_Address + Block_Base_Address_Offset

To calculate the bank number and the block number from the block base address:

If the address is in the range of the parameter bank, the Bank Number is 0 and the Block Number can be directly read from Tables *36* and *42* (parameter bank Block Addresses), in the Block Number column, in the row that corresponds to the address given. Otherwise, the Block Number can be calculated using the formulas below:

For the top configuration (M58LR256KT and M58LR128KT):

 $Block_Number = ((NOT address) / 2^{16}) + 3$

For the bottom configuration (M58LR256KB and M58LR128KB):

Block_Number = (address / 2¹⁶) + 3

For both configurations the Bank Number and the Block Number Offset can be calculated using the following formulas:

 $Bank_Number = (Block_Number - 3) / 8$

Block_Number_Offset = Block_Number - 3 - (Bank_Number x 8)

Block number	Size (Kwords)	Address range
0	16	7FC000-7FFFFF
1	16	7F8000-7FBFFF
2	16	7F4000-7F7FFF
3	16	7F0000-7F3FFF
4	64	7E0000-7EFFFF
5	64	7D0000-7DFFFF
6	64	7C0000-7CFFFF
7	64	7B0000-7BFFFF
8	64	7A0000-7AFFFF
9	64	790000-79FFFF
10	64	780000-78FFFF

 Table 33.
 M58LR128KT - Parameter bank block addresses

Table 34. M58LR128KT - main bank base addresses

Bank number ⁽¹⁾	Block numbers	Bank base address
1	11-18	70 0000
2	19-26	68 0000
3	27-34	60 0000
4	35-42	58 0000
5	43-50	50 0000
6	51-58	48 0000
7	59-66	40 0000
8	67-74	38 0000
9	75-82	30 0000
10	83-90	28 0000
11	91-98	20 0000
12	99-106	18 0000
13	107-114	10 0000
14	115-122	08 0000
15	123-130	00 0000

1. There are two Bank Regions: Bank Region 1 contains all the banks that are made up of main blocks only; Bank Region 2 contains the banks that are made up of the parameter and main blocks (parameter bank).

Block base address offset				
07 0000				
06 0000				
05 0000				
04 0000				
03 0000				
02 0000				
01 0000				
00 0000				

 Table 35.
 M58LR128KT - block addresses in main banks

Table 36. M58LR256KT - parameter bank block addresses

Block number	Size (Kwords)	Address range		
0	16	FFC000-FFFFFF		
1	16	FF8000-FFBFFF		
2	16	FF4000-FF7FFF		
3	16	FF0000-FF3FFF		
4	64	FE0000-FEFFFF		
5	64	FD0000-FDFFFF		
6	64	FC0000-FCFFFF		
7	64	FB0000-FBFFFF		
8	64	FA0000-FAFFFF		
9	64	F90000-F9FFFF		
10	64	F80000-F8FFFF		
11	1 64 F70000-F7FFF			
12	64 F60000-F6FFF			
13	64	F50000-F5FFFF		
14	64	F40000-F4FFFF		
15	64	F30000-F3FFFF		
16	64	F20000-F2FFFF		
17	64	F10000-F1FFFF		
18	64	F00000-F0FFFF		



Bank number ⁽¹⁾	Block numbers	Bank base address	
1	19-34	E00000	
2	35-50	D00000	
3	51-66	C00000	
4	67-82	B00000	
5	83-98	A00000	
6	99-114	900000	
7	115-130	800000	
8	131-146	700000	
9	147-162	600000	
10	163-178	500000	
11	179-194	400000	
12	195-210	300000	
13	211-226	200000	
14	227-242	100000	
15	243-258	000000	

Table 37. M58LR256KT - main bank base addresses

1. There are two Bank Regions: Bank Region 1 contains all the banks that are made up of main blocks only; Bank Region 2 contains the banks that are made up of the parameter and main blocks (parameter bank).

Table 38. M58LR256KT - block addresses in main banks

Block number offset	Block base address offset
0	0F0000
1	0E0000
2	0D0000
3	0C0000
4	0B0000
5	0A0000
6	090000
7	080000
8	070000
9	060000
10	050000
11	040000
12	030000
13	020000
14	010000
15	000000



Block number	Size (Kwords)	Address range	
10 64		070000-07FFFF	
9	64	060000-06FFFF	
8	64	050000-05FFFF	
7	64	040000-04FFFF	
6	64	030000-03FFFF	
5	64	020000-02FFFF	
4 64		010000-01FFFF	
3	16	00C000-00FFFF	
2	16	008000-00BFFF	
1	16	004000-007FFF	
0	16	000000-003FFF	

 Table 39.
 M58LR128KB - parameter bank block addresses

Table 40. M58LR128KB - main bank base addresses

Bank number ⁽¹⁾ Block numbers		Bank base address	
15 123-130		78 0000	
14	115-122	70 0000	
13	107-114	68 0000	
12	99-106	60 0000	
11	91-98	58 0000	
10	83-90	50 0000	
9	75-82	48 0000	
8 67-74		40 0000	
7	59-66	38 0000	
6	51-58	30 0000	
5	43-50	28 0000	
4	35-42	20 0000	
3	27-34	18 0000	
2	19-26	10 0000	
1 11-18		08 0000	

 There are two Bank Regions: Bank Region 2 contains all the banks that are made up of main blocks only; Bank Region 1 contains the banks that are made up of the parameter and main blocks (parameter bank).

Block number offset	Block base address offset	
7	070000	
6	060000	
5	050000	
4	040000	
3	030000	
2	020000	
1	010000	
0	000000	

Table 41. M58LR128KB - block addresses in main banks

Table 42. M58LR256KB - parameter bank block addresses

Block number	Size (Kwords)	Address range	
18	64	0F0000-0FFFF	
17	64	0E0000-0EFFFF	
16	64	0D0000-0DFFFF	
15	64	0C0000-0CFFFF	
14	64	0B0000-0BFFFF	
13	64	0A0000-0AFFFF	
12	64	090000-09FFFF	
11	64	080000-08FFFF	
10	64	070000-07FFFF	
9	64	060000-06FFFF	
8	64	050000-05FFFF	
7	64	040000-04FFFF	
6	64	030000-03FFFF	
5	64	020000-02FFFF	
4	64	010000-01FFFF	
3	16	00C000-00FFFF	
2	16	008000-00BFFF	
1	16	004000-007FFF	
0	16	000000-003FFF	

Bank number	Block numbers	Bank base address		
15	243-258	F00000		
14	227-242	E00000		
13	211-226	D00000		
12	195-210	C00000		
11	179-194	B00000		
10	163-178	A00000		
9	147-162	900000		
8	131-146	800000		
7	115-130	700000		
6	99-114 600000			
5	83-98	500000		
4	67-82	400000		
3	51-66	300000		
2	35-50	200000		
1 19-34		100000		

Table 43. M58LR256KB - main bank base addresses

1. There are two Bank Regions: Bank Region 2 contains all the banks that are made up of main blocks only; Bank Region 1 contains the banks that are made up of the parameter and main blocks (parameter bank).

Block number offset	Block base address offset
15	0F0000
14	0E0000
13	0D0000
12	0C0000
11	0B0000
10	0A0000
9	090000
8	080000
7	070000
6	060000
5	050000
4	040000
3	030000
2	020000
1	010000
0	000000

Appendix B Common Flash interface

The common Flash interface is a JEDEC approved, standardized data structure that can be read from the Flash memory device. It allows a system software to query the device to determine various electrical and timing parameters, density information and functions supported by the memory. The system can interface easily with the device, enabling the software to upgrade itself when necessary.

When the read CFI Query command is issued the device enters CFI query mode and the data structure is read from the memory. Tables *45*, *46*, *47*, *48*, *49*, *50*, *51*, *52*, *53* and *54* show the addresses used to retrieve the data. The query data is always presented on the lowest order data outputs (DQ0-DQ7), the other outputs (DQ8-DQ15) are set to 0.

The CFI data structure also contains a security area where a 64 bit unique security number is written (see *Figure 6: Protection Register memory map*). This area can be accessed only in Read mode by the final user. It is impossible to change the security number after it has been written by Numonyx. Issue a Read Array command to return to read mode.

Offset	Sub-section name	Description	
000h	Reserved	Reserved for algorithm-specific information	
010h	CFI query identification string	Command set ID and algorithm data offset	
01Bh	System interface information	Device timing and voltage information	
027h	Device geometry definition	Flash device layout	
Р	Primary algorithm-specific extended query table	Additional information specific to the primary algorithm (optional)	
A	Alternate algorithm-specific extended query table	Additional information specific to the alternate algorithm (optional)	
080h	Security code area	Lock Protection Register Unique device number and User programmable OTP	

Table 45. Query structure overview

1. The Flash memory display the CFI data structure when CFI Query command is issued. In this table are listed the main sub-sections detailed in Tables *46*, *47*, *48* and *49*. Query data is always presented on the lowest order data outputs.

Offset	Sub-section Name	Description	Value	
000h	0020h	Manufacturer code		Numonyx
	88C4h		M58LR128KT	Тор
0016	88C5h	Device and	M58LR128KB	Bottom
001h	880Dh	Device code	M58LR256KT	Тор
	880Eh		M58LR256KB	Bottom
002h-00Fh	Reserved	Reserved		
010h	0051h	Query Unique ASCII String "QRY"		"Q"
011h	0052h			"R"
012h	0059h			"Y"
013h	0001h	Primary Algorithm Command Set and Control Interface ID code 16 bit ID code defining a specific algorithm		
014h	0000h			
015h	offset = P = 000Ah	Address for Primary Algorithm extended Query table		n 104h
016h	0001h	(see Table 49)		p = 10Ah
017h	0000h	Alternate Vendor Command Set and Control		
018h	0000h	Interface ID Code second vendor - specified algorithm supported		NA
019h	value = A = 0000h	Address for Alternate Algorithm extended Query table		NA
01Ah	0000h			INA

 Table 46.
 CFI query identification string



Offset	Data	Description	Value
01Bh	0017h	V _{DD} Logic Supply Minimum Program/Erase or Write voltage bit 7 to 4 BCD value in volts bit 3 to 0 BCD value in 100 millivolts	1.7V
01Ch	0020h	V _{DD} Logic Supply Maximum Program/Erase or Write voltage bit 7 to 4 BCD value in volts bit 3 to 0 BCD value in 100 millivolts	2V
01Dh	0085h	V _{PP} [Programming] Supply Minimum Program/Erase voltage bit 7 to 4 HEX value in volts bit 3 to 0 BCD value in 100 millivolts	8.5V
01Eh	0095h	V _{PP} [Programming] Supply Maximum Program/Erase voltage bit 7 to 4 HEX value in volts bit 3 to 0 BCD value in 100 millivolts	9.5V
01Fh	0004h	Typical time-out per single byte/word program = $2^n \mu s$	16µs
020h	0009h	Typical time-out for Buffer Program = $2^n \mu s$	512µs
021h	000Ah	Typical time-out per individual block erase = 2 ⁿ ms	1s
022h	0000h	Typical time-out for full chip erase = 2 ⁿ ms	NA
023h	0004h	Maximum time-out for word program = 2 ⁿ times typical	256µs
024h	0004h	Maximum time-out for Buffer Program = 2 ⁿ times typical	8192µs
025h	0002h	Maximum time-out per individual block erase = 2 ⁿ times typical	4s
026h	0000h	Maximum time-out for chip erase = 2 ⁿ times typical	NA

Table 47. CFI query system interface information

Tak	Die 48.	Device	jeometry definition	
C	Offset	Data	Description	Value
	0018 027h		M58LR128KT/B device size = 2 ⁿ in number of bytes	16 Mbytes
	02711	0019h	M58LR256KT/B device size = 2 ⁿ in number of bytes	32 Mbytes
	028h 029h	0001h 0000h	Flash Device Interface Code description	x16 Async.
	02Ah 02Bh	0006h 0000h	Maximum number of bytes in multi-byte program or page = 2 ⁿ	64 bytes
(02Ch	0002h	Number of identical sized erase block regions within the device bit 7 to 0 = x = number of Erase Block Regions	2
	02Dh	007Eh 0000h	M58LR128KT/B Erase Block Region 1 Information Number of identical-size erase blocks = 007Eh+1	127
	02Eh	00FEh 0000h	M58LR256KT/B Erase Block Region 1 Information Number of identical-size erase blocks = 00FEh+1	255
DEVICES	02Fh 030h	0000h 0002h	Erase Block Region 1 Information Block size in Region 1 = 0200h * 256 byte	128 Kbyte
TOP DE	031h 032h	0003h 0000h	Erase Block Region 2 Information Number of identical-size erase blocks = 0003h+1	4
	033h 034h	0080h 0000h	Erase Block Region 2 Information Block size in Region 2 = 0080h * 256 byte	32 Kbyte
	035h 038h	Reserved	Reserved for future erase block region information	NA
	02Dh 02Eh	0003h 0000h	Erase Block Region 1 Information Number of identical-size erase block = 0003h+1	4
S	02Fh 030h	0080h 0000h	Erase Block Region 1 Information Block size in Region 1 = 0080h * 256 bytes	32 Kbytes
OM DEVICES	031h	007Eh 0000h	M58LR128KT/B Erase Block Region 2 Information Number of identical-size erase block = 007Eh+1	127
BOTTOM	032h	00FEh 0000h	M58LR256KT/B Erase Block Region 2 Information Number of identical-size erase block = 00FEh+1	255
BC	033h 034h	0000h 0002h	Erase Block Region 2 Information Block size in Region 2 = 0200h * 256 bytes	128 Kbytes
	035h 038h	Reserved	Reserved for future erase block region information	NA

Table 48. Device geometry definition



Offset	Data	Description	Value	
(P)h = 10Ah	0050h		"P"	
	0052h	Primary Algorithm extended Query table unique ASCII string "PRI"	"R"	
	0049h		" "	
(P+3)h =10Dh	0031h	Major version number, ASCII	"1"	
(P+4)h = 10Eh	0033h	Minor version number, ASCII	"3"	
(P+5)h = 10Fh	00E6h	Extended Query table contents for Primary Algorithm. Address		
	0003h	(P+5)h contains less significant byte.		
(P+7)h = 111h (P+8)h = 112h	0000h	bit 0 Chip Erase supported(1 = Yes, 0 = No) bit 1 Erase Suspend supported(1 = Yes, 0 = No) bit 2 Program Suspend supported(1 = Yes, 0 = No) bit 3 Legacy Lock/Unlock supported(1 = Yes, 0 = No) bit 4 Queued Erase supported(1 = Yes, 0 = No) bit 5 Instant individual block locking supported(1 = Yes, 0 = No) bit 6 Protection bits supported(1 = Yes, 0 = No) bit 7 Page mode read supported(1 = Yes, 0 = No) bit 8 Synchronous read supported(1 = Yes, 0 = No) bit 9 Simultaneous operation supported(1 = Yes, 0 = No) bit 10 to 31 Reserved; undefined bits are '0'. If bit 31 is '1' then another 31 bit field of optional features follows at the end of the bit-30 field.	No Yes No No Yes Yes Yes Yes	
(P+9)h = 113h	0001h	Supported Functions after Suspend Read Array, Read Status Register and CFI Query bit 0 Program supported after Erase Suspend (1 = Yes, 0 = No) bit 7 to 1 Reserved; undefined bits are '0'		
(P+A)h = 114h	(P+A)h = 114h 0003h Block Protect Status Defines which bits in the Block Status Register section of the Query are implemented.			
(P+B)h = 115h	0000h	bit 0 Block protect Status Register Lock/Unlock bit active (1 = Yes, 0 = No) bit 1 Block Lock Status Register Lock-Down bit active (1 = Yes, 0 = No) bit 15 to 2 Reserved for future use; undefined bits are '0'	Yes Yes	

 Table 49.
 Primary algorithm-specific extended query table

Offset	Data	Description			
(P+C)h = 116h	0018h	V _{DD} Logic Supply Optimum Program/Erase voltage (highest performance) bit 7 to 4 HEX value in volts bit 3 to 0 BCD value in 100 mV	1.8V		
(P+D)h = 117h 0090h		V _{PP} Supply Optimum Program/Erase voltage bit 7 to 4 HEX value in volts bit 3 to 0 BCD value in 100 mV	9V		

 Table 49.
 Primary algorithm-specific extended query table

Table 50. Protection Register information

Offset	Data	Description	Value			
(P+E)h = 118h	0002h	Number of protection register fields in JEDEC ID space. 0000h indicates that 256 fields are available.	2			
(P+F)h = 119h	0080h	Protection Field 1: Protection Description	80h			
(P+10)h = 11Ah	0000h	Bits 0-7 Lower byte of protection register address	00h			
(P+ 11)h = 11Bh	0003h	Bits 8-15 Upper byte of protection register address Bits 16-23 2 ⁿ bytes in factory pre-programmed region	8 bytes			
(P+12)h = 11Ch	0003h	Bits 24-31 2 ⁿ bytes in user programmable region				
(P+13)h = 11Dh	0089h	Protocilian Devictor & Protocilian Developintion	89h			
(P+14)h = 11Eh	0000h	Protection Register 2: Protection Description Bits 0-31 protection register address Bits 32-39 n number of factory programmed regions (lower				
(P+15)h = 11Fh	0000h					
(P+16)h = 120h	0000h	byte)				
(P+17)h = 121h	0000h	Bits 40-47 n number of factory programmed regions (upper byte)				
(P+18)h = 122h	0000h	Bits 48-55 2 ⁿ bytes in factory programmable region	0			
(P+19)h = 123h	0000h	Bits 56-63 n number of user programmable regions (lower byte) Bits 64-71 n number of user programmable regions (upper byte) Bits 72-79 2 ⁿ bytes in user programmable region				
(P+1A)h = 124h	0010h					
(P+1B)h = 125h	0000h					
(P+1C)h = 126h	0004h		16			

Table 51.Burst read information

Offset	Data	Description			
(P+1D)h = 127h	0003h	Page-mode read capability bits 0-7 n' such that 2 ⁿ HEX value represents the number of read-page bytes. See offset 0028h for device word width to determine page-mode data output width.			
(P+1E)h = 128h	0004h	Number of synchronous mode read configuration fields that follow.	4		

Table 51. Burst read information	Table 51.	Burst read information
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Offset	Data	Description			
(P+1F)h = 129h	0001h	Synchronous mode read capability configuration 1 bit 3-7 Reserved bit 0-2 n' such that 2 ⁿ⁺¹ HEX value represents the maximum number of continuous synchronous reads when the device is configured for its maximum word width. A value of 07h indicates that the device is capable of continuous linear bursts that will output data until the internal burst counter reaches the end of the device's burstable address space. This field's 3-bit value can be written directly to the read configuration register bit 0-2 if the device is configured for its maximum word width. See offset 0028h for word width to determine the burst data output width.	4		
(P+20)h = 12Ah	0002h	Synchronous mode read capability configuration 2			
(P-21)h = 12Bh	0003h	Synchronous mode read capability configuration 3	16		
(P+22)h = 12Ch	0007h	007h Synchronous mode read capability configuration 4			

Table 52. Bank and erase block region information^{(1) (2)}

Flash memory (top)		(top) Flash memory (bottom)		Description
Offset	Data	Offset Data		Description
(P+23)h = 12Dh	02h	(P+23)h = 12Dh	02h	Number of Bank Regions within the device

1. The variable P is a pointer that is defined at CFI offset 015h.

2. Bank Regions. There are two Bank Regions, see *Table 34*, *Table 37*, *Table 40* and *Table 43*.

Table 53. Bank and erase block region 1 information

Flash memory (top)		Flash memory (bottom)		Description
Offset ⁽¹⁾	Data	Offset ⁽¹⁾	Data	Description
(P+24)h = 12Eh	0Fh	(P+24)h = 12Eh	01h	Number of identical banks within Bank Region
(P+25)h = 12Fh	00h	(P+25)h = 12Fh	00h	1
(P+26)h = 130h	11h	(P+26)h = 130h	11h	Number of program or erase operations allowed in Bank Region 1: Bits 0-3: Number of simultaneous program operations Bits 4-7: Number of simultaneous erase operations
(P+27)h = 131h	00h	(P+27)h = 131h	00h	Number of program or erase operations allowed in other banks while a bank in same region is programming Bits 0-3: Number of simultaneous program operations Bits 4-7: Number of simultaneous erase operations

Flash memory	(top)	Flash memory (b	pottom)	Description
Offset ⁽¹⁾	Data	Offset ⁽¹⁾	Data	Description
(P+28)h = 132h	00h	(P+28)h = 132h	00h	Number of program or erase operations allowed in other banks while a bank in this region is erasing Bits 0-3: Number of simultaneous program operations Bits 4-7: Number of simultaneous erase operations
(P+29)h = 133h	01h	(P+29)h = 133h	02h	Types of erase block regions in Bank Region 1 n = number of erase block regions with contiguous same-size erase blocks. Symmetrically blocked banks have one blocking region ⁽²⁾
(P+2A)h = 134h	07h ⁽³⁾ 0Fh ⁽⁴⁾	(P+2A)h = 134h	03h	Bank Region 1 Erase Block Type 1 Information Bits 0-15: n+1 = number of identical-sized
(P+2B)h = 135h	00h	(P+2B)h = 135h	00h	erase blocks
(P+2C)h = 136h	00h	(P+2C)h = 136h	80h	Bits 16-31: n x 256 = number of bytes in erase block region
(P+2D)h = 137h	02h	(P+2D)h = 137h	00h	DIOCK TEGION
(P+2E)h = 138h	64h	(P+2E)h = 138h	64h	Bank Region 1 (Erase Block Type 1)
(P+2F)h = 139h	00h	(P+2F)h = 139h	00h	Minimum block erase cycles x 1000
(P+30)h = 13Ah	01h	(P+30)h = 13Ah	01h	Bank Region 1 (Erase Block Type 1): Blts per cell, internal ECC Bits 0-3: bits per cell in erase region Bit 4: reserved for "internal ECC used" Blts 5-7: reserved
(P+31)h = 13Bh	03h	(P+31)h = 13Bh	03h	Bank Region 1 (Erase Block Type 1): Page mode and Synchronous mode capabilities Bit 0: Page-mode reads permitted ⁽⁵⁾ Bit 1: Synchronous reads permitted Bit 2: Synchronous writes permitted Bits 3-7: reserved
		(P+32)h = 13Ch	06h ⁽³⁾ 0Eh ⁽⁴⁾	Bank Region 1 Erase Block Type 2 Information
		(P+33)h = 13Dh	00h	Bits 0-15: n+1 = number of identical-sized erase blocks
		(P+34)h = 13Eh	00h	Bits 16-31: n x 256 = number of bytes in erase
		(P+35)h = 13Fh	02h	block region
		(P+36)h = 140h	64h	Bank Region 1 (Erase Block Type 2)
		(P+37)h = 141h	00h	Minimum block erase cycles x 1000

 Table 53.
 Bank and erase block region 1 information (continued)



Flash memory (top)		Flash memory (b	pottom)	
Offset ⁽¹⁾	Data	Offset ⁽¹⁾	Data	Description
		(P+38)h = 142h	01h	Bank Regions 1 (Erase Block Type 2): Blts per cell, internal ECC Bits 0-3: bits per cell in erase region Bit 4: reserved for "internal ECC used" Blts 5-7: reserved
		(P+39)h = 143h	03h	Bank Region 1 (Erase Block Type 2): Page mode and Synchronous mode capabilities Bit 0: Page-mode reads permitted Bit 1: Synchronous reads permitted Bit 2: Synchronous writes permitted Bits 3-7: reserved

Table 53.	Bank and erase block region 1 information	(continued)
Table 53.	bank and erase block region 1 information	(continuea)

1. The variable P is a pointer that is defined at CFI offset 015h.

2. Bank Regions. There are two Bank Regions, see Table 34, Table 37, Table 40 and Table 43.

3. Applies to M58LR128KT/B only.

4. Applies to M58LR256KT/B only.

5. Although the device supports Page Read mode, this is not described in the datasheet as its use is not advantageous in a multiplexed device.

Table 54.	Bank and erase block region 2 information
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Flash memory (top)		Flash memory (bottom)		Description	
Offset ⁽¹⁾	Data	Offset ⁽¹⁾	Data	Description	
(P+32)h = 13Ch	01h	(P+3A)h = 144h	0Fh	Number of identical banks within Bank Region	
(P+33)h = 13Dh	00h	(P+3B)h = 145h	00h	2	
(P+34)h = 13Eh	11h	(P+3C)h = 146h	11h	Number of program or erase operations allowed in Bank Region 2: Bits 0-3: Number of simultaneous program operations Bits 4-7: Number of simultaneous erase operations	
(P+35)h = 13Fh	00h	(P+3D)h = 147h	00h	Number of program or erase operations allowed in other banks while a bank in this region is programming Bits 0-3: Number of simultaneous program operations Bits 4-7: Number of simultaneous erase operations	
(P+36)h = 140h	00h	(P+3E)h = 148h	00h	Number of program or erase operations allowed in other banks while a bank in this region is erasing Bits 0-3: Number of simultaneous program operations Bits 4-7: Number of simultaneous erase operations	



Flash memory (top)		Flash memory (bottom)		
Offset ⁽¹⁾	Data	Offset ⁽¹⁾	Data	Description
(P+37)h = 141h	02h	(P+3F)h = 149h	01h	Types of erase block regions in Bank Region 2 n = number of erase block regions with contiguous same-size erase blocks. Symmetrically blocked banks have one blocking region. ⁽²⁾
(P+38)h = 142h	06h ⁽³⁾ 0Eh ⁽⁴⁾	(P+40)h = 14Ah	07h ⁽³⁾ 0Fh ⁽⁴⁾	Bank Region 2 Erase Block Type 1 Information Bits 0-15: n+1 = number of identical-sized
(P+39)h = 143h	00h	(P+41)h = 14Bh	00h	erase blocks
(P+3A)h = 144h	00h	(P+42)h = 14Ch	00h	Bits 16-31: n x 256 = number of bytes in erase
(P+3B)h = 145h	02h	(P+43)h = 14Dh	02h	block region
(P+3C)h = 146h	64h	(P+44)h = 14Eh	64h	Bank Region 2 (Erase Block Type 1)
(P+3D)h = 147h	00h	(P+45)h = 14Fh	00h	Minimum block erase cycles x 1000
(P+3E)h = 148h	01h	(P+46)h = 150h	01h	Bank Region 2 (Erase Block Type 1): Blts per cell, internal ECC Bits 0-3: bits per cell in erase region Bit 4: reserved for "internal ECC used" Blts 5-7: reserved
(P+3F)h = 149h	03h	(P+47)h = 151h	03h	Bank Region 2 (Erase Block Type 1):Page mode and Synchronous mode capabilities (defined in <i>Table 51</i>) Bit 0: Page-mode reads permitted ⁽⁵⁾ Bit 1: Synchronous reads permitted Bit 2: Synchronous writes permitted Bits 3-7: reserved
(P+40)h = 14Ah	03h			Bank Region 2 Erase Block Type 2 Information
(P+41)h = 14Bh	00h			Bits 0-15: n+1 = number of identical-sized erase blocks
(P+42)h = 14Ch	80h			Bits 16-31: n x 256 = number of bytes in erase
(P+43)h = 14Dh	00h			block region
(P+44)h = 14Eh	64h			Bank Region 2 (Erase Block Type 2)
(P+45)h = 14Fh	00h			Minimum block erase cycles x 1000
(P+46)h = 150h	01h			Bank Region 2 (Erase Block Type 2): Blts per cell, internal ECC Bits 0-3: bits per cell in erase region Bit 4: reserved for "internal ECC used" Blts 5-7: reserved

 Table 54.
 Bank and erase block region 2 information (continued)



Flash memory (top)		Flash memory (bottom)		Description
Offset ⁽¹⁾	Data	Offset ⁽¹⁾ Data		Description
(P+47)h = 151h	03h			Bank Region 2 (Erase Block Type 2): Page mode and Synchronous mode capabilities (defined in <i>Table 51</i>) Bit 0: Page-mode reads permitted Bit 1: Synchronous reads permitted Bit 2: Synchronous writes permitted Bits 3-7: reserved
(P+48)h = 152h		(P+48)h = 152h		Feature Space definitions
(P+49)h = 153h		(P+43)h = 153h		Reserved

 Table 54.
 Bank and erase block region 2 information (continued)

1. The variable P is a pointer which is defined at CFI offset 015h.

2. Bank Regions. There are two bank regions, see Table 34, Table 37, Table 40 and Table 43.

3. Applies to M58LR128KT/B only.

4. Applies to M58LR256KT/B only.

5. Although the device supports Page Read mode, this is not described in the datasheet as its use is not advantageous in a multiplexed device.

Appendix C Flowcharts and pseudocodes



Figure 24. Program flowchart and pseudocode

 Status check of SR1 (Protected Block), SR3 (V_{PP} Invalid) and SR4 (Program Error) can be made after each program operation or after a sequence.

2. If an error is found, the Status Register must be cleared before further Program/Erase Controller operations.

3. Any address within the bank can equally be used.





1. Any address within the bank can equally be used.

2. If an error is found, the Status Register must be cleared before further Program/Erase operations.



Figure 26. Buffer program flowchart and pseudocode

1. n + 1 is the number of data being programmed.

2. Next Program data is an element belonging to buffer_Program[].data; Next Program address is an element belonging to buffer_Program[].address

3. Routine for Error Check by reading SR3, SR4 and SR1.





Figure 27. Program suspend and resume flowchart and pseudocode

1. The Read Status Register command (Write 70h) can be issued just before or just after the Program Resume command.



Figure 28. Block erase flowchart and pseudocode

1. If an error is found, the Status Register must be cleared before further Program/Erase operations.

2. Any address within the bank can equally be used.





Figure 29. Erase suspend and resume flowchart and pseudocode

1. The Read Status Register command (Write 70h) can be issued just before or just after the Erase Resume command.



Figure 30. Locking operations flowchart and pseudocode

1. Any address within the bank can equally be used.




Figure 31. Protection Register program flowchart and pseudocode

1. Status check of SR1 (Protected Block), SR3 (V_{PP} Invalid) and SR4 (Program Error) can be made after each program operation or after a sequence.

2. If an error is found, the Status Register must be cleared before further Program/Erase Controller operations.

3. Any address within the bank can equally be used.



Figure 32. Buffer enhanced factory program flowchart and pseudocode

N numonyx

Appendix D Command interface state tables

							Com	mand Input						
Current	CI State	Read Array ⁽²⁾ (FFh)	Program Setup ⁽³⁾⁽⁴⁾ (10/40h)	Buffer Program ⁽³⁾⁽⁴⁾ (E8h)	Block Erase, Setup ⁽³⁾⁽⁴⁾ (20h)	BEFP Setup (80h)	Blank Check setup (BCh)	Erase Confirm P/E Resume, Block Unlock confirm, BEFP Confirm ⁽³⁾⁽⁴⁾ (D0h)	Blank Check confirm (CBh)	Buffer Program, Program/ Erase Suspend (B0h)	Read Status Register (70h)	Clear Status Register ⁽⁵⁾ (50h)	Read Electronic Signature , Read CFI Query (90h, 98h)	
Rea	ady	Ready	Program Setup	BP Setup	Erase Setup	BEFP Setup	Blank Check setup			Read	Register (70h) Hegister (5) , Real CFI Qui (90h, 9i dy (Lock Error)			
Lock/Cl	R Setup			Ready (Lo	ck Error)			Ready (unlock block)		Rea	dy (Lock E	Error)		
	Setup						С	TP Busy						
ОТР	Busy	OTP Busy	IS in OTP Busy	OTP busy	IS in OTF	P Busy			(OTP Busy				
	IS in OTP busy						C	TP Busy						
	Setup						Pro	gram Busy						
	Busy	Program Busy	IS in Program Busy	Program Busy	IS in Pro Bus			Program Busy	Program Suspend	Program Busy				
Program	IS in Program Busy	Program Busy								ady eady (Lock Error) n d Program Busy rogram Suspend d 2 (data load) ad 2 the first address) Ready (error) Buffar Program Busy				
	Suspend	PS	IS in PS	PS	IS in Pro Suspe		PS	Program Busy		Pro	gram Sus	oend		
	IS in PS						Progr	am Suspend						
	Setup				Buffe	er Progra	am Load 1	l (give word cou	int load (N	N-1));				
	Buffer Load 1		if	N=0 go to	Buffer Prog	ıram Cor	nfirm. Else	e (N ≠ 0) go to E	Buffer Pro	gram Load 2	2 (data loa	d)		
	Buffer Load 2		(note					count =0; Else I any block addre				ress)		
Buffer	Confirm			Ready ((error)			BP Busy		F	leady (erro	or)		
Program	Busy	BP Busy	IS in BP Busy	BP Busy	IS in BP	Busy		BP Busy		BP Suspend	Buffe	er Progran	n Busy	
	IS in BP Busy						Buffer	Program Busy						
	Suspend	BP Suspend	IS in BP Suspend	BP Suspend	IS in BP S	uspend	BP Suspend	BP busy		Buffer	Program S	Suspend		
	IS in BP Suspend						Buffer Pr	ogram Suspend	b					

Table 55.	Command interface states - modify table, next state ⁽¹⁾



							Com	mand Input	,		,					
Current	CI State	Read Array ⁽²⁾ (FFh)	Program Setup ⁽³⁾⁽⁴⁾ (10/40h)	Buffer Program (3)(4) (E8h)	Block Erase, Setup ⁽³⁾⁽⁴⁾ (20h)	BEFP Setup (80h)	Blank Check setup (BCh)	Erase Confirm P/E Resume, Block Unlock confirm, BEFP Confirm ⁽³⁾⁽⁴⁾ (D0h)	Blank Check confirm (CBh)	Buffer Program, Program/ Erase Suspend (B0h)	Read Status Register (70h)	Clear Status Register ⁽⁵⁾ (50h)	Read Electronic Signature , Read CFI Query (90h, 98h)			
	Setup			Ready	(error)			Erase Busy		F	leady (erro	or)				
	Busy	Erase Busy	IS in Erase Busy	Erase Busy	IS in Eras	e Busy		Erase Busy		Erase Suspend						
Erase	IS in Erase Busy						Er	Erase Busy								
	Suspend	Erase Suspend	Program in ES	BP in ES	IS in El Suspe		ES	Erase Busy		Er	ase Susp	Ise Suspend Program Busy in Erase Suspend				
	IS in ES						Eras	se Suspend				se Suspend Program Busy in Erase				
	Setup					Pro	ogram Bus	sy in Erase Susp	pend							
BusyProgram Busy in ESIS in Program Busy in ESProgram Busy in ESIS in Program Busy in ESProgram Busy in ESProgram Program Busy in ES																
Program in Erase Suspend	IS in Program busy in ES	Program busy in Erase Suspend														
	Suspend	PS in ES	IS in PS in ES	PS in ES	IS in Pro Suspend		PS in ES	Program Busy in ES	P	Program Sus	pend in E	Erase Suspend				
	IS in PS in ES	Program Suspend in Erase Suspend														
	Setup	Buffer P	rogram Loa	d 1 in Era	se Suspend	l (give w		load (N-1)); if N Program Load 2	=0 go to I	Buffer Progr	am confiri	n. Else (N	≠ 0) go to			
	Buffer Load 1				Buffe	er Progra	am Load 2	in Erase Suspe	end (data	load)						
	Buffer Load 2	Buffer Pro	ogram Conf					lse Buffer Progr ddress is differe				note: Buff	er Program			
	Confirm		Erase	Suspend (sequence e	rror)		BP Busy in ES		Erase Sus	pend (seq	uence erro	or)			
Buffer Program in Erase	Busy	BP Busy in ES	IS in BP Busy in ES	BP busy in ES	IS in BP t ES			BP Busy in ES		BP Suspend in ES	Buffer F	Program B	usy in ES			
Suspend	IS in BP busy in ES					Buffer	Program	Busy in Erase S	Suspend							
	Suspend	BP Suspend in ES	IS in BP Suspend in ES	BP Suspend in ES	IS in BP S in Erase S	uspend uspend	BP Suspend in ES	BP Busy in Erase Suspend	Buffe	er Program	Suspend i	n Erase S	uspend			
	IS in BP Suspend in ES			•	•	В	P Suspend	d in Erase Susp	end							

Table 55. Command interface states - modify table, next state ⁽¹⁾
--

							Com	mand Input					
Current	CI State	Read Array ⁽²⁾ (FFh)	Program Setup ⁽³⁾⁽⁴⁾ (10/40h)	Buffer Program ⁽³⁾⁽⁴⁾ (E8h)	Block Erase, Setup ⁽³⁾⁽⁴⁾ (20h)	BEFP Setup (80h)	Blank Check setup (BCh)	Erase Confirm P/E Resume, Block Unlock confirm, BEFP Confirm ⁽³⁾⁽⁴⁾ (D0h)	Blank Check confirm (CBh)	Buffer Program, Program/ Erase Suspend (B0h)	Read Status Register (70h)	Clear Status Register (5) (50h)	Read Electronic Signature , Read CFI Query (90h, 98h)
Blank	Setup		Ready (error)						Blank Check busy		Ready	(error)	
Check	Busy	Blank Check busy	IS in Blank Check busy		Blank Check busy								
	Setup in suspend		Eras	e Suspend	d (Lock Erro	or)		Erase Suspend		Erase S	uspend (Lo	ock Error)	
Buffer	Setup			Ready ((error)			BEFP Busy		F	Ready (erro	or)	
EFP	Busy						BE	FP Busy ⁽⁶⁾					

Table 55.	Command interface states - modify table, next state ⁽¹⁾ (continued)
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1. CI = Command Interface, CR = Configuration register, BEFP = Buffer Enhanced Factory program, P/E C = Program/Erase controller, IS = Illegal State, BP = Buffer Program, ES = Erase Suspend.

2. At power-up, all banks are in read array mode. Issuing a Read Array command to a busy bank, results in undetermined data output.

3. The two cycle command should be issued to the same bank address.

4. If the P/E C is active, both cycles are ignored.

5. The Clear Status Register command clears the SR error bits except when the P/E C. is busy or suspended.

6. BEFP is allowed only when Status Register bit SR0 is reset to '0'. BEFP is busy if Block Address is first BEFP Address. Any other commands are treated as data.



Table 56. CC						-	Command Input	-				
Current CI State	Read Array ⁽³⁾ (FFh)	Program Setup ⁽⁴⁾ (5) (10/40h)	Buffer Program (E8h)	Block Erase, Setup ⁽⁴⁾ ⁽⁵⁾ (20h)	BEFP Setup (80h)	Blank Check setup (BCh)	Erase Confirm P/E Resume, Block Unlock confirm, BEFP Confirm ⁽⁴⁾⁽⁵⁾ (D0h)	Blank Check confir m (CBh)	Program/ Erase Suspend (B0h)	Read Status Register (70h)	Clear Status Register (50h)	Read Electronic signature, Read CFI Query (90h, 98h)
Program Setup		I.										
Erase Setup												
OTP Setup												
Program Setup in Erase Suspend												
BEFP Setup												
BEFP Busy												
Buffer Program Setup												
Buffer Program Load 1												
Buffer Program Load 2												
Buffer Program Confirm							Status Register					
Buffer Program Setup in Erase Suspend												
Buffer Program Load 1 in Erase Suspend												
Buffer Program Load 2 in Erase Suspend												
Buffer Program Confirm in Erase Suspend												
Blank Check setup												
Lock/CR Setup												
Lock/CR Setup in Erase Suspend												

Table 56. Command interface states - modify table, next output state^{(1) (2)}



							Command Input	t				
Current CI State	Read Array ⁽³⁾ (FFh)	Program Setup ⁽⁴⁾ (5) (10/40h)	Buffer Program (E8h)	Block Erase, Setup ⁽⁴⁾ ⁽⁵⁾ (20h)	BEFP Setup (80h)	Blank Check setup (BCh)	Erase Confirm P/E Resume, Block Unlock confirm, BEFP Confirm ⁽⁴⁾⁽⁵⁾ (D0h)	Blank Check confir m (CBh)	Program/ Erase Suspend (B0h)	Read Status Register (70h)	Clear Status Register (50h)	Read Electronic signature, Read CFI Query (90h, 98h)
OTP Busy												Status Register
Ready												
Program Busy												
Erase Busy												
Buffer Program Busy												
Program/Erase Suspend												
Buffer Program Suspend	Array		Statu	s Registe	r		Output	Unchang	ged	Status Register	Output Unchanged	Electronic
Program Busy in Erase Suspend												Signature/ CFI
Buffer Program Busy in Erase Suspend												
Program Suspend in Erase Suspend												
Buffer Program Suspend in Erase Suspend												
Blank Check busy												
Illegal State						C	Output Unchange	d				

Table 56.	Command interface states - modif	y table, next out	put state ^{(1) (2)}	(continued)
Table 56.	Command interface states - modif	y table, next out	put state ^{(1) (2)} ((contin	ued

- The output state shows the type of data that appears at the outputs if the bank address is the same as the command address. A bank can be placed in Read Array, Read Status Register, Read Electronic Signature or Read CFI mode, depending on the command issued. Each bank remains in its last output state until a new command is issued to that bank. The next state does not depend on the bank output state.
- 2. CI = Command Interface, CR = Configuration Register, BEFP = Buffer Enhanced Factory Program, P/E. C. = Program/Erase Controller.
- 3. At Power-Up, all banks are in read array mode. Issuing a Read Array command to a busy bank, results in undetermined data output.
- 4. The two cycle command should be issued to the same bank address.
- 5. If the P/EC is active, both cycles are ignored.



					Comman	d Input								
Curre	nt CI State	Lock/CR Setup ⁽²⁾ (60h)	OTP Setup ⁽²⁾ (C0h)	Block Lock Confirm (01h)	Block Lock- Down Confirm (2Fh)	Set CR Confirm (03h)	Block Address (WA0) ⁽³⁾ (XXXXh)	lllegal Command ⁽⁴⁾	P/E C operation completed (5)					
F	Ready	Lock/CR Setup	OTP Setup			N/A								
Lock/	CR Setup	Ready (Lock	error)		Ready		Ready	(Lock error)	N/A					
	Setup			(OTP Busy				N/A					
ОТР	Busy	IS in OTP Busy OTP Busy							Ready					
	IS in OTP busy OTP Busy							IS Ready						
	Setup	Program Busy					N/A							
	Busy	IS in Program	Busy			Program Bu	ısy	Ready						
Program	IS in Program busy	m Program busy							Program busy					IS Ready
	Suspend	IS in PS Program Suspend												
	IS in PS		Program Suspend				N/A							
	Setup		Buffer Pi	rogram Load	1 (give word	count load (N	-1));		N/A					
	Buffer Load 1		Buffer Program Load 2 ⁽⁶⁾ Exit s		see note ⁽⁶⁾	N/A								
	Buffer Load 2	Buffer Program Cor t	nfirm when co his point if an					ogram will fail at	N/A					
	Confirm			Re	eady (error)				N/A					
Buffer Program	Busy	IS in BP Bu	isy		В	uffer Progran	n Busy		Ready					
	IS in Buffer Program busy			Buffer	Program Bu	sy			IS Ready					
	Suspend	IS in BP Sus	pend		Buf	fer Program S	Suspend							
	IS in BP Suspend			Buffer P	rogram Susp	end			N/A					
	Setup Ready (error)								N/A					
	Busy	IS in Erase E	Busy			Erase Bus	Sy .		Ready					
Erase	IS in Erase busy			E	IS ready									
	Suspend	Lock/CR Setup in ES	IS in ES			Erase Susp	end		N/A					
	IS in ES			Era	se Suspend				11/7					

Table 57	Command interface states - lock table, next state ⁽¹⁾
Table 57.	Command Interface states - lock table, next state

					Comman	d Input					
Currei	nt CI State	Lock/CR Setup ⁽²⁾ (60h)	OTP Setup ⁽²⁾ (C0h)	Block Lock Confirm (01h)	Block Lock- Down Confirm (2Fh)	Set CR Confirm (03h)	Block Address (WA0) ⁽³⁾ (XXXXh)	lllegal Command ⁽⁴⁾	P/E C operation completed (5)		
	Setup			Program Bu	sy in Erase S	Suspend			N/A		
	Busy	IS in Program bu	sy in ES		Program	n Busy in Era	se Suspend		ES		
Program in Erase Suspend	IS in Program busy in ES			Program Bu	Program Busy in Erase Suspend						
	Suspend	IS in PS in ES Program Suspend in Erase Suspend									
	IS in PS in ES		Program Suspend in Erase Suspend								
	Setup	Buffer Program Load 1 in Erase Suspend (give word count load (N-1))									
	Buffer Load 1	Buffer	Program Loa	ad 2 in Erase							
	Buffer Load 2	Buffer Program C Suspend (note: E				N/A					
Buffer	Confirm	Erase Suspend (sequence error)									
Program in Erase Suspend	Busy	IS in BP busy	in ES		Buffer Prog	gram Busy in	Erase Suspe	nd	ES		
Suspenu	IS in BP busy in ES			BP	IS in ES						
	Suspend	IS in BP suspen	d in ES		Buffer Progra	am Suspend i	in Erase Susp	pend			
	IS in BP Suspend in ES		Buff	er Program S	uspend in Er	ase Suspend			N/A		
Blank	Setup			Re	eady (error)				N/A		
Check	Blank Check busy	IS in Blank Che	ck busy		I	Blank Check	busy		Ready		
Lock/CR	Setup in ES	Erase Suspend (L	ock error)	E	rase Suspen	d	Erase Susp	N/A			
BEED	Setup			Re	eady (error)		•		N/A		
BEFP	Busy		BEF	P Busy ⁽⁸⁾			Exit	BEFP Busy ⁽⁸⁾	N/A		

	Table 57.	Command interface states	 lock table, next state⁽¹⁾ 	(continued)
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 CI = Command Interface, CR = Configuration register, BEFP = Buffer Enhanced Factory program, P/E C = Program/Erase controller, IS = Illegal State, BP = Buffer program, ES = Erase suspend, WA0 = Address in a block different from first BEFP address.

- 2. If the P/E C is active, both cycle are ignored.
- 3. BEFP Exit when Block Address is different from first Block Address and data are FFFFh.
- 4. Illegal commands are those not defined in the command set.
- 5. N/A: not available. In this case the state remains unchanged.
- 6. If N=0 go to Buffer Program Confirm. Else (not =0) go to Buffer Program Load 2 (data load)
- 7. If N=0 go to Buffer Program Confirm in Erase suspend. Else (not =0) go to Buffer Program Load 2 in Erase suspend.
- 8. BEFP is allowed only when Status Register bit SR0 is set to '0'. BEFP is busy if Block Address is first BEFP Address. Any other commands are treated as data.



					Comma	and Input				
Current CI State	Lock/CR Setup ⁽³⁾ (60h)	Blank Check setup (BCh)	OTP Setup ⁽³⁾ (C0h)	Blank Check confirm (CBh)	Block Lock Confirm (01h)	Block Lock- Down Confirm (2Fh)	Set CR Confirm (03h)	BEFP Exit ⁽⁴⁾ (FFFFh)	Illegal Command ⁽⁵⁾	P. E./C. Operation Completed
Program Setup										
Erase Setup										
OTP Setup										
Program Setup in Erase Suspend										
BEFP Setup										
BEFP Busy										
Buffer Program Setup										
Buffer Program Load 1										
Buffer Program Load 2					Status Regis	ter				
Buffer Program Confirm										Output Unchanged
Buffer Program Setup in Erase Suspend										
Buffer Program Load 1 in Erase Suspend										
Buffer Program Load 2 in Erase Suspend										
Buffer Program Confirm in Erase Suspend										
Blank Check setup]									
Lock/CR Setup										
Lock/CR Setup in Erase Suspend			State	us Register			Array	Status	Register	

Table 58. Command interface states - lock table, next output state ^{(1) (2)}



			Command Input							
Current CI State	Lock/CR Setup ⁽³⁾ (60h)	Blank Check setup (BCh)	OTP Setup ⁽³⁾ (C0h)	Blank Check confirm (CBh)	Block Lock Confirm (01h)	Block Lock- Down Confirm (2Fh)	Set CR Confirm (03h)	BEFP Exit ⁽⁴⁾ (FFFFh)	Illegal Command ⁽⁵⁾	P. E./C. Operation Completed
OTP Busy										
Ready	Status Register							Array	Output Unchanged	
Program Busy										
Erase Busy										
Buffer Program Busy				Output Unchanged						
Program/Erase Suspend										
Buffer Program Suspend										
Program Busy in Erase Suspend										
Buffer Program Busy in Erase Suspend										
Program Suspend in Erase Suspend										
Buffer Program Suspend in Erase Suspend										
Blank Check busy										
Illegal State				Output Unchanged						

Table 58. Command interface states - lock table, next output state (continued)^{(1) (2)}

- The output state shows the type of data that appears at the outputs if the bank address is the same as the command address. A bank can be placed in Read Array, Read Status Register, Read Electronic Signature or Read CFI mode, depending on the command issued. Each bank remains in its last output state until a new command is issued to that bank. The next state does not depend on the bank's output state.
- 2. CI = Command Interface, CR = Configuration Register, BEFP = Buffer Enhanced Factory Program, P/E. C. = Program/Erase Controller.
- 3. If the P/EC is active, both cycles are ignored.
- 4. BEFP Exit when Block Address is different from first Block Address and data are FFFFh.
- 5. Illegal commands are those not defined in the command set.



15 Revision history

Table 59. Document revision histor	/
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Date	Revision	Changes
23-Mar-2007	1	Initial release.
		In <i>Table 18: Program/erase times and endurance cycles</i> , changed the suspend latency program and erase values to 5, 10 and 5, 20 respectively, to 20 and 25 for both. In <i>Table 19: Absolute maximum ratings</i> , changed the V _{IO} maximum value from 3.8 to V _{DDQ} + 0.6 V. In <i>Table 22: DC characteristics - currents</i> :
06-Dec-2007	2	 Added the 256 Mbit values for I_{DD2}, I_{DD3}, and I_{DD7}. For I_{DD5}, changed the supply current (both program and erase) V_{PPH} values to 10 and 30, and changed the V_{DD} and values to 20 and 35, respectively.
		 For I_{DD6}, changed the "asynchronous read in another bank" values to 33 and 50, respectively; changed the "synchronous read in another bank" values to 50 and 67, respectively. In <i>Table 28: Reset and power-up AC characteristics</i>, changed the t_{VDHPH} value from 250 to 300.
27-Mar-2008	3	Applied Numonyx branding.
01-Jul-2008	4	Added Figure 3: TFBGA88 connections (top view through package), Figure 21: TFBGA88 8 x 10 mm - 8 x 10 ball array, 0.8 mm pitch, bottom view package outline, Table 23: VFBGA79 9x11mm - 0.75 pitch -package outline, and amended Table 32: Ordering information scheme to provide the ZB package information, which is available for the M58LR128KT/B devices.
11-Sep-2008	5	Add 79B vfBGA 9x11 package for 256Mbit. Add 88b package 8x10. Change from T=-25°C to T=-30°C Change Terase for Boot sectors (0.4s to 0.6s)
23-Oct-2008	6	Change the package reference from ZA to ZQ.

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